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The CP/M Operating System Manual was prepared using the Digital Research TEX Text Formatter and printed in the United States of America.

*****
* First Edition: 1976 *
* Second Edition: July 1982 *
* Third Edition: March 1983 *
* Fourth Edition: March 1984 *
*****



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[illegible]



## Section 1

### CP/M Features and Facilities

#### 1.1 Introduction

CP/M is a monitor control program for microcomputer system development that uses floppy disks or Winchester hard disks for backup storage. Using a computer system based on the Intel 8080 microcomputer, CP/M provides an environment for program construction, storage, and editing, along with assembly and program check-out facilities. CP/M can be easily altered to execute with any computer configuration that uses a Zilog Z80 or an Intel 8080 Central Processing Unit (CPU) and has at least 20K bytes of main memory with up to 16 disk drives. A detailed discussion of the modifications required for any particular hardware environment is given in Section 6. Although the standard Digital Research version operates on a single-density Intel Model 800, microcomputer development system several different hardware manufacturers support their own input-output (I/O) drivers for CP/M.

The CP/M monitor provides rapid access to programs through a comprehensive file management package. The file subsystem supports a named file structure, allowing dynamic allocation of file space as well as sequential and random file access. Using this file system, CP/M can be easily altered to execute with any computer configuration that uses a Zilog Z80 or an Intel 8080 Central Processing Unit (CPU) and has at least 20K bytes of main memory with up to 16 disk drives. A detailed discussion of the modifications required for any particular hardware environment is given in Section 6. Although the standard Digital Research version operates on a single-density Intel Model 800, microcomputer development system several different hardware manufacturers support their own input-output (I/O) drivers for CP/M.

CP/M 2 is a high-performance, single console operating system that uses table-driven techniques to allow field reconfiguration to match a wide variety of disk capacities. All fundamental file restrictions are removed, maintaining upward compatibility from previous versions of release 1.

Features of CP/M 2 include field specification of one to sixteen logical drives, each containing up to eight megabytes. Any particular file can reach the full drive size with the capability of expanding to thirty-two megabytes in future releases. The directory size can be field-configured to contain any reasonable number of entries, and each file is optionally tagged with Read-Only and system attributes. Users of CP/M 2 are physically separated by user numbers, with facilities for file copy operations from one user area to another. Powerful relative-record random access functions are present in CP/M 2 that provide direct access to any of the 65536 records of an eight-megabyte file.

CP/M also supports ED, a powerful context editor, ASM, an Intel-compatible assembler, and DDT, debugger subsystems. Optional software includes a powerful Intel-compatible macro assembler, symbolic debugger, along with various high-level languages. When coupled with CP/M's Console Command Processor (CCP), the resulting facilities equal or exceed similar large computer facilities.



CP/M is logically divided into several distinct parts:

- o BIOS (Basic I/O System), hardware-dependent
- o BDOS (Basic Disk Operating System)
- o CCP (Console Command Processor)
- o TPA (Transient Program Area)

The BIOS provides the primitive operations necessary to access the disk drives and to interface standard peripherals: teletype, CRT, paper tape reader/punch, and user-defined peripherals. You can tailor peripherals for any particular hardware environment by patching this portion of CP/M. The BDOS provides disk management by controlling one or more disk drives containing independent file directories. The BDOS implements disk allocation strategies that provide fully dynamic file construction while minimizing head movement across the disk during access. The BDOS has entry points that include the following primitive operations, which the program accesses:

- o SEARCH looks for a particular disk file by name.
- o OPEN opens a file for further operations.
- o CLOSE closes a file after processing.
- o RENAME changes the name of a particular file.
- o READ reads a record from a particular file.
- o WRITE writes a record to a particular file.
- o SELECT selects a particular disk drive for further operations.

the remainder of the CP/M system. The CCP reads the console device and processes commands, which include listing the file directory, printing the contents of files, and controlling the operation of transient programs, such as assemblers, editors, and debuggers. The standard commands that are available in the CCP are listed in Section 1.2.1.

The last segment of CP/M is the area called the Transient Program Area (TPA). The TPA holds programs that are loaded from the disk under command of the CCP. During program editing, for example, the TPA holds the CP/M text editor machine code and data areas. Similarly, programs created under CP/M can be checked out by loading and executing these programs in the TPA.

Any or all of the CP/M component subsystems can be overlaid by an executing program. That is, once a user's program is loaded into the TPA, the CCP, BDOS, and BIOS areas can be used as the program's data area. A bootstrap loader is programmatically accessible whenever the BIOS portion is not overlaid; thus, the user program need only branch to the bootstrap loader at the end of execution and the complete CP/M monitor is reloaded from disk.

The CP/M operating system is partitioned into distinct modules, including the BIOS portion that defines the hardware environment in which CP/M is executing. Thus, the standard system is easily modified to any nonstandard environment by changing the peripheral drivers to handle the custom system.



[illegible]



## 1.2 Functional Description



[illegible]



```
b>A: Switch back to A.
```

afn	ambiguous file reference
-----	--------------------------

1	2	3	4	5	6	7
---	---	---	---	---	---	---

1-6



CP/M Operating System Manual 1.4 Built-in Commands

ERA \*.\* Erase all files on the current disk. In this case, the CCP prompts the console with the message

ALL FILES (Y/N)?

which requires a Y response before files are actually removed.

ERA b:\*.PRN All files on drive B that satisfy the ambiguous reference ???????.PRN are deleted, independently of the currently logged disk.

#### 1.4.2 DIR Command

Syntax:

DIR afn

The DIR (directory) command causes the names of all files that satisfy the ambiguous filename afn to be listed at the console device. As a special case, the command

DIR

lists the files on the currently logged disk (the command DIR is equivalent to the command DIR \*.\*). The following are valid DIR commands:

DIR X.Y  
DIR X?Z.C?M  
DIR ???.Y

Similar to other CCP commands, the afn can be preceded by a drive name. The following DIR commands cause the selected drive to be addressed before the directory search takes place:

DIR B:  
DIR B:X.Y  
DIR B:\*.A?M

If no files on the selected disk satisfy the directory request, the message NO FILE appears at the console.



CP/M Operating System Manual	1.4 Built-in Commands
------------------------------	-----------------------

### 1.4.3 REN Command

Syntax:

[illegible]

The REN (rename) command allows you to change the names of files on disk. The file satisfying ufn2 is changed to ufn1. The currently logged disk is assumed to contain the file to rename (ufn2). You can also type a left-directed arrow instead of the equal sign if the console supports this graphic character. The following are examples of the REN command:

REN	X.Y=O.R	The file O.R is changed to X.Y
-----	---------	--------------------------------

[illegible]

The operator precedes either ufn1 or ufn2 (or both) by an optional drive address. If ufn1 is preceded by a drive name, then ufn2 is assumed to exist on the same drive. Similarly, if ufn2 is preceded by a drive name, then ufn1 is assumed to exist on the drive as well. The same drive must be specified in both cases if both ufn1 and ufn2 are preceded by drive names. The following REN commands illustrate this format:

REN A:X.ASM=Y.ASM	The file Y.ASM is changed to X.ASM on drive A.
-------------------	---

REN B:ZAP.BAS=ZOT.BAS	The file ZOT.BAS is changed to ZAP.BAS on drive B.
-----------------------	--

REN B:A.ASM=B:A.BAK	The file A.BAK is renamed to A.ASM on drive B.
---------------------	---

```

    If ufn1 is already present, the REN command responds with the
error FILE EXISTS and not perform the change.  If ufn2 does not
exist on the specified disk, the message NO FILE is printed at the
console.

```

#### 1.4.4 SAVE Command

Syntax:
---------

[illegible]

The SAVE command places n pages (256-byte blocks) onto disk from the TPA and names this file ufn. In the CP/M distribution system, the TPA starts at 100H (hexadecimal) which is the second page of memory. The SAVE command must specify 2 pages of memory if

[illegible]



[illegible]



another logical area within the same directory. Drives that are logged-in while addressing one user number are automatically active when the operator moves to another. A user number is simply a prefix that accesses particular directory entries on the active disks.

The active user number is maintained until changed by a subsequent USER command, or until a cold start when user 0 is again assumed.

### 1.5 Line Editing and Output Control

The CCP allows certain line-editing functions while typing command lines. The CTRL-key sequences are obtained by pressing the control and letter keys simultaneously. Further, CCP command lines are generally up to 255 characters in length; they are not acted upon until the carriage return key is pressed.

Table 1-1. Line-editing Control Characters

Character	Meaning
CTRL-C	Reboots CP/M system when pressed at start of line.
CTRL-D	Ends current input; if no character has been typed, but line is not sent until the carriage return key is pressed.
CTRL-H	Backspaces one character position.
CTRL-J	Terminates current input (line feed).
CTRL-M	Terminates current input (carriage return).
CTRL-P	Copies all subsequent console output to the currently assigned list device (see Section 1.6.1). Output is sent to the list device and the console device until the next CTRL-P is pressed.
CTRL-R	Retypes current command line; types a clean line following character deletion with rubouts.
CTRL-S	Stops the console output temporarily. Program execution and output continue when you press any character at the console, for example another CTRL-S. This feature stops output on high speed consoles, such as CRTs, in order to view a segment of output before continuing.



Table 1-1. (continued)

Character	Meaning
CTRL-U	Deletes the entire line typed at the console.
CTRL-X	Same as CTRL-U.
CTRL-Z	Ends input from the console (used in PIP and ED).
RUB/DEL	Deletes and echoes the last character typed at the console.

1.6 Transient Commands

Transient commands are loaded from the currently logged disk and executed in the TPA. The transient commands for execution under the CCP are below. Additional functions are easily defined by the user (see Section 1.6.3).

Table 1-2. CP/M Transient Commands

Command	Function
STAT	Lists the number of bytes of storage remaining on the currently logged disk, provides statistical information about particular files, and displays or alters device assignment.
ASM	Loads the CP/M assembler and assembles the specified program from disk.
LOAD	Loads the file in Intel HEX machine code format and produces a file in machine executable form which can be loaded into the TPA. This loaded program becomes a new command under the CCP.
DDT	Loads the CP/M debugger into TPA and starts execution.
PIP	Loads the Peripheral Interchange Program for subsequent disk file and peripheral transfer operations.
ED	Loads and executes the CP/M text editor program.
SYSGEN	Creates a new CP/M system disk.















CP/M Operating System Manual	1.6 Transient Commands
------------------------------	------------------------

It is emphasized that the physical device names might not actually correspond to devices that the names imply. That is, you can implement the PTP: device as a cassette write operation. The exact correspondence and driving subroutine is defined in the BIOS portion of CP/M. In the standard distribution version of CP/M, these devices correspond to their names on the Model 800 development system.

The command,

[illegible]

```
produces a summary of the available status commands, resulting in
the output:
```

Temp R/O Disk	d:\$R/O
Set Indicator:	filename.type \$R/O \$R/W \$SYS \$DIR
Disk Status:	DSK: d:DSK
Iobyte Assign:	

which gives an instant summary of the possible STAT commands and shows the permissible logical-to-physical device assignments:

				CON:	=	TTY:	CRT:	BAT:	UC1:
				RDR:	=	TTY:	PTR:	UR1:	UR2:
				PUN:	=	TTY:	PTP:	UP1:	UP2:

The logical device to the left takes any of the four physical assignments shown to the right. The current logical-to-physical mapping is displayed by typing the command:

[illegible]

This command produces a list of each logical device to the left and the current corresponding physical device to the right.. For example, the list might appear as follows:

				CON:	=	CRT:
				RDR:	=	UR1:
				PUN:	=	PTP:
				LST:	=	TTY:

The current logical-to-physical device assignment is changed by typing a STAT command of the form:

STAT	ld1 = pd1, ld2 = pd2, ..., ldn = pdn
------	--------------------------------------

where `ld1` through `ldn` are logical device names and `pd1` through `pdn` are compatible physical device names. For example, `ldi` and `pdj` appear on the same line in the `VAL:` command shown above. The following example shows valid `STAT` commands that change the current logical-to-physical device assignments:



```
STAT CON:=CRT:
```

```
STAT PUN:=TTY:, LST:=LPT:, RDR:=TTY:
```

The command form,

```
STAT d:filename.typ $S
```

where d: is an optional drive name and filename.typ is an unambiguous or ambiguous filename, produces the following output display format:

Size	Recs	Bytes	Ext	Acc
------	------	-------	-----	-----

48	48	6K	1 R/O	A:ED.COM
----	----	----	-------	----------

55	55	12K	1 R/O	(A:PIP.COM)
----	----	-----	-------	-------------

65536	128	16K	2 R/W	A:X.DAT
-------	-----	-----	-------	---------

where the \$S parameter causes the Size field to be displayed. Without the \$S, the Size field is skipped, but the remaining fields are displayed. The Size field lists the virtual file size in records, while the Recs field sums the number of virtual records in each extent. For files constructed sequentially, the Size and Recs fields are identical. The Bytes field lists the actual number of bytes allocated to the corresponding file. The minimum allocation unit is determined at configuration time; thus, the number of bytes

the last allocated block for sequential files. Random access files are given data areas only when written, so the Bytes field contains the only accurate allocation figure. In the case of random access, the Size field gives the logical end-of-file record position and the Recs field counts the logical records of each extent. Each of these extents, however, can contain unallocated holes even though they are added into the record count.

The Ext field counts the number of physical extents allocated to the file. The Ext count corresponds to the number of directory entries given to the file. Depending on allocation size, there can be up to 128K bytes (8 logical extents) directly addressed by a single directory entry. In a special case, there are actually 256K bytes that can be directly addressed by a physical extent.

The Acc field gives the R/O or R/W file indicator, which you can change using the commands shown. The four command forms,

```
STAT d:filename.typ $R/O
```

```
STAT d:filename.typ $R/W
```

```
STAT d:filename.typ $SYS
```

```
STAT d:filename.typ $DIR
```

set or reset various permanent file indicators. The R/O indicator places the file, or set of files, in a Read-Only status until changed by a subsequent STAT command. The R/O status is recorded in the directory with the file so that it remains R/O through



CP/M Operating System Manual 1.6 Transient Commands

intervening cold start operations. The R/W indicator places the file in a permanent Read-Write status. The SYS indicator attaches the system indicator to the file, while the DIR command removes the system indicator. The filename type may be ambiguous or unambiguous, but files whose attributes are changed are listed at the console when the change occurs. The drive name denoted by d: is optional.

When a file is marked R/O, subsequent attempts to erase or write into the file produce the following BDOS message at your screen:

```
BDOS Err on d: File R/O
```

```
lists the drive characteristics of the disk named by d: that is in
the range A:., B:.,...,P:.. The drive characteristics are listed in
the following format:
```

	d:	Drive Characteristics	
65536:	128	Byte Record Capacity	
8192:		Kilobyte Drive Capacity	
128:	32	Byte Directory Entries	
0:		Checked Directory Entries	
1024:		Records/Extent	
128:		Records/Block	
58:		Sectors/Track	
2:		Reserved Tracks	

capacity (65536 is an eight-megabyte drive), followed by the total capacity listed in kilobytes. The directory size is listed next, followed by the checked entries. The number of checked entries is usually identical to the directory size for removable media, because this mechanism is used to detect changed media during CP/M operation without an intervening warm start. For fixed media, the number is usually zero, because the media are not changed without at least a cold or warm start.

The number of records per extent determines the addressing capacity of each directory entry (1024 times 128 bytes, or 128K in the previous example). The number of records per block shows the basic allocation size (in the example, 128 records/block times 128 bytes per record, or 16K bytes per block). The listing is then followed by the number of physical sectors per track and the number of reserved tracks.

For logical drives that share the same physical disk, the number of reserved tracks can be quite large because this mechanism is used to skip lower-numbered disk areas allocated to other logical disks. The command form

STAT	DSK:
------	------

produces a drive characteristics table for all currently active drives. The final STAT command form is



[illegible]



is also produced, which contains 8080 machine language in Intel HEX format suitable for subsequent loading and execution (see Section 1.6.3). For complete details of CP/M's assembly language program, see Section 3.

[illegible]

loads the assembler from the currently logged drive and processes the source program ALPHA.ASM on drive B. The HEX and PRN files are also placed on drive B in this case.

### Syntax:

LOAD ufn

The `LOAD` command reads the file `ufn`, which is assumed to contain HEX format machine code, and produces a memory image file

be of the form:

					X	.	H	E	X
--	--	--	--	--	---	---	---	---	---

and only the filename X need be specified in the command. The LOAD command creates a file named

					X.COM
--	--	--	--	--	-------

that marks it as containing machine executable code. The file is actually loaded into memory and executed when the user types the filename X immediately after the prompting character > printed by the CCP.

Generally, the CCP reads the filename X following the prompting character and looks for a built-in function name. If no function name is found, the CCP searches the system disk directory for a file by the name

					X.COM
--	--	--	--	--	-------

If found, the machine code is loaded into the TPA, and the program executes. Thus, the user need only LOAD a hex file once; it can be subsequently executed any number of times by typing the primary name. This way, you can invent new commands in the CCP. Initialized disks contain the transient commands as COM files, which are optionally deleted. The operation takes place on an alternate drive if the filename is prefixed by a drive name. Thus,



CP/M Operating System Manual 1.6 Transient Commands

[illegible]

```
brings the LOAD program into the TPA from the currently logged disk
and operates on drive B after execution begins.
```

**Note:** the BETA.HEX file must contain valid Intel format hexadecimal machine code records (as produced by the ASM program, for example) that begin at 100H of the TPA. The addresses in the hex records must be in ascending order; gaps in unfilled memory regions are filled with zeroes by the LOAD command as the hex records are read. Thus, LOAD must be used only for creating CP/M standard COM files that operate in the TPA. Programs that occupy regions of memory other than the TPA are loaded under DDT.

1	.	6	.	4					P	I	P
---	---	---	---	---	--	--	--	--	---	---	---

Syntax:
---------

[illegible]

```
PIP destination=source#1, source#2, ..., source #n
```

PIP is the CP/M Peripheral Interchange Program that implements the basic media conversion operations necessary to load, print, punch, copy, and combine disk files. The PIP program is initiated by typing one of the following forms:

[illegible]

PIP command line

In both cases PIP is loaded into the MPA and executed. In the first form, PIP reads command lines directly from the console, prompted with the \* character, until an empty command line is typed (for example, a single carriage return is issued by the operator). Each successive command line causes some media conversion to take place according to the rules shown below.

In the second form, the PIP command is equivalent to the first, except that the single command line given with the PIP command is automatically executed, and PIP terminates immediately with no further prompting of the console for input command lines. The form of each command line is

```
destination = source#1, source#2, ..., source#n
```

where destination is the file or peripheral device to receive the data, and source#1, ..., source#n is a series of one or more files or devices that are copied from left to right to the destination.

When multiple files are given in the command line (for example, n>1), the individual files are assumed to contain ASCII characters, with an assumed CP/M end-of-file character (CTRL-Z) at the end of each file (see the O parameter to override this assumption). Lower-case ASCII alphabets are internally translated to upper-case to be



The destination and source elements are unambiguous references to CP/M source files with or without a preceding disk drive name. That is, any file can be referenced with a preceding drive name (A: through P:) that defines the particular drive where the file can be obtained or stored. When the drive name is not included, the currently logged disk is assumed. The destination file can also appear as one or more of the source files, in which case the source file is not altered until the entire concatenation is complete. If it already exists, the destination file is removed if the command line is properly formed. It is not removed if an error condition arises. The following command lines, with explanations to the right, are valid as input to PIP:

```
Copies to file X from file Y, where
X and Y are unambiguous filenames;
Y remains unchanged.
```

```
Concatenates files Y and Z and
copies to file X, with Y and Z
unchanged.
```

```
concatenation of the Y and Z.ASM
files.
```

```
Moves a copy of OLD.ZAPP from drive
B to the currently logged disk;
names the file NEW.ZOT.
```

```
Concatenates file B.V from drive B
with C.W from drive a and D.X from
the logged disk; creates the file
A.U on drive b.
```

For convenience, PIP allows abbreviated commands for transferring files between disk drives. The abbreviated PIP forms are

PIP	$d := afn$							
PIP	$d_1 := d$	$d := afn$						
PIP	$ufn \equiv d$	$d :=$						
PIP	$d_1 := ufn$	$d \equiv d_2$						

The first form copies all files from the currently logged disk that satisfy the afn to the same files on drive d, where  $d = A \dots P$ . The second form is equivalent to the first, where the source for the copy is drive d<sub>1</sub>, where  $d_1 = A \dots P$ . The third form is equivalent to the command `PIP d1:ufn=d2:ufn` which copies the file given by ufn



1	-	2	2
---	---	---	---



1	-	2	3
---	---	---	---







You can also specify one or more PIP parameters, enclosed in left and right square brackets, separated by zero or more blanks. Each parameter affects the copy operation, and the enclosed list of parameters must immediately follow the affected file or device. Generally, each parameter can be followed by an optional decimal integer value (the S and Q parameters are exceptions). Table 1-4 describes valid PIP parameters.

Table 1-4. PIP Parameters

Parameter	Meaning
B	Blocks mode transfer. Data are buffered by PIP until an ASCII x-off character, CTRL-S, is received from the source device. This allows transfer of data to a disk file from a continuous reading device, such as a cassette reader. Upon receipt of the x-off, PIP clears the disk buffers and returns for more input data. The amount of data that can be buffered depends on the memory size of the host system. PIP issues an error message if the buffers overflow.
Dn	Deletes characters that extend past column n from the character source. This parameter is generally used to truncate long lines that are sent to a narrow printer or console device.
E	Echoes all transfer operations to the console as they are being performed.
F	Filters form-feeds from the file. All embedded form-feeds are removed. The P parameter can be used simultaneously to insert new form-feeds.
Gn	Gets file from user number n (n in the range 0-15).
H	Transfers HEX data. All data are checked for proper Intel hex file format. Nonessential characters between hex records are removed during the copy operation. The console is prompted for corrective action in case errors occur.
I	Ignores :00 records in the transfer of Intel hex format file. The I parameter automatically sets the H parameter.



## Transient Commands

Table 1-4. (continued)

Parameter	Meaning
L	Translates upper-case alphabetic to lower-case.
N	Adds line numbers to each line transferred to the destination, starting at one and incrementing by 1. Leading zeroes are suppressed, and the number is followed by a colon. If N2 is specified, leading zeroes are included and a tab is inserted following the number. The tab is expanded if T is set.
O	Transfers non-ASCII object files. The normal CP/M end-of-file is ignored.
Pn	Includes page ejects at every n lines with an initial page eject. If n = 1 or is excluded altogether, page ejects occur every 60 lines. If the F parameter is used, form-feed suppression takes place before the new page ejects are inserted.
Qs^Z	Quits copying from the source device or file when the string s, terminated by CTRL-Z, is

R	Reads system files.
Ss^Z	Start copying from the source device when the string s, terminated by CTRL-Z, is encountered. The S and Q parameters can be used to abstract a particular section of a file, such as a subroutine. The start and quit strings are always included in the copy operation.
	If you specify a command line after the PIP command keyword, the CCP translates strings following the S and Q parameters to upper-case. If you do not specify a command line, PIP does not perform the automatic upper-case translation.
Tn	Expands tabs, CTRL-I characters, to every nth column during the transfer of characters to the destination from the source.
U	Translates lower-case alphabetic to upper-case during the copy operation.



Meaning
---------

Verifies that data have been copied correctly by rereading after the write operation (the destination must be a disk file).

Writes over R/O files without console interrogation.	
--	--

Zeros the parity bit on input for each ASCII character.	
---	--

P	I	P	X	.	A	S	M	=	B	:	[	v	]
---	---	---	---	---	---	---	---	---	---	---	---	---	---

Copies X.ASM from drive B to the current drive and verifies that the data were properly copied.

```
PIP LPT:=X.ASM[nt8u]
```

Copies X.ASM to the LPT: device; numbers each line, expands tabs to every eighth column, and translates lower-case alphabetic to upper-case.

```
PIP  PUN:=X.HEX[i],Y.ZOT[h]
```

First copies X.HEX to the PUN: device and ignores the trailing :00 record in X.HEX; continues the transfer of data by reading Y.ZOT, which contains HEX records, including any :00 records it contains.

```
PIP X.LIB=Y.ASM[sSUBRI:^z qJMP L3^z]
```

Copies from the file Y.ASM into the file X.LIB. The command starts the copy when the string SUBR1: has been found, and quits copying after the string JMP L3 is encountered.



1	-	2	8
---	---	---	---



[illegible]











```
SOURCE ON d THEN TYPE RETURN
```

Place a disk containing the CP/M operating system on drive d (d is one of A, B, C, or D). Answer by pressing a carriage return when ready.

FUNCTION COMPLETE

System is copied to memory. SYSGEN then prompts with the following:

DESTINATION DRIVE NAME  
(OR RETURN TO REBOOT)

If a disk is being initialized, place the new disk into a drive and answer with the drive name. Otherwise, press a carriage return and the system reboots from drive A. Typing drive name d causes SYSGEN to prompt with the following message:

```
DESTINATION ON d
THEN TYPE RETURN
```

```
Place new disk into drive d; press return
when ready.
```

FUNCTION COMPLETE

```
New disk is initialized in drive d.
```

The DESTINATION prompt is repeated until a single carriage return is pressed at the console, so that more than one disk can be initialized.

Upon completion of a successful system generation, the new disk contains the operating system, and only the built-in commands are available. An IBM-compatible disk appears to CP/M as a disk with an empty directory; therefore, the operator must copy the appropriate COM files from an existing CP/M disk to the newly constructed disk using the PIP transient.

You can copy all files from an existing disk by typing the following PIP command:

P	I	P	B:	=	A:	*	.	*	[	v	]
---	---	---	----	---	----	---	---	---	---	---	---



This command copies all files from disk drive A to disk drive B and verifies that each file has been copied correctly. The name of each file is displayed at the console as the copy operation proceeds.

Note that a SYSGEN does not destroy the files that already exist on a disk; it only constructs a new operating system. If a disk is being used only on drives B through P and will never be the source of a bootstrap operation on drive A, the SYSGEN need not take place.

#### 1.6.7 SUBMIT Command

##### Syntax:

```
SUBMIT ufn parm#1 .... parm#n
```

The SUBMIT command allows CP/M commands to be batched for automatic processing. The ufn given in the SUBMIT command must be the filename of a file that exists on the currently logged disk, with an assumed file type of SUB. The SUB file contains CP/M prototype commands with possible parameter substitution. The actual parameters parm#1 .... parm#n are substituted into the prototype commands, and, if no errors occur, the file of substituted commands are processed sequentially by CP/M.

The prototype command file is created using the ED program,

```
$1 $2 $3 .... $n
```

corresponding to the number of actual parameters that will be included when the file is submitted for execution. When the SUBMIT transient is executed, the actual parameters parm#1 .... parm#n are paired with the formal parameters \$1 .... \$n in the prototype commands. If the numbers of formal and actual parameters do not correspond, the SUBMIT function is aborted with an error message at the console. The SUBMIT function creates a file of substituted commands with the name

```
$$$ .SUB
```

on the logged disk. When the system reboots, at the termination of the SUBMIT, this command file is read by the CCP as a source of input rather than the console. If the SUBMIT function is performed on any disk other than drive A, the commands are not processed until the disk is inserted into drive A and the system reboots. You can abort command processing at any time by pressing the rubout key when the command is read and echoed. In this case, the \$\$\$ .SUB file is removed and the subsequent commands come from the console. Command processing is also aborted if the CCP detects an error in any of the commands. Programs that execute under CP/M can abort processing of command files when error conditions occur by erasing any existing \$\$\$ .SUB file.



## 1.6 Transient Commands

The last command in a SUB file can initiate another SUB file, allowing chained batch commands:

ASM	\$1
DIR	\$1.*
ERA	*.BAK
PIP	\$2:= \$1.PRN
ERA	\$1.PRN

SUBMIT ASMBL X PRN

[illegible]

An additional utility program called XSUB extends the power of the SUBMIT facility to include line input to programs as well as the CCP. The XSUB command is included as the first line of the SUBMIT file. When it is executed, XSUB self-relocates directly below the CCP. All subsequent SUBMIT command lines are processed by XSUB so that programs that read buffered console input, BDOS Function 10, receive their input directly from the SUBMIT file. For example, the file SAVER.SUB can contain the following SUBMIT lines:

XSUB
DDT
\$1.COM
R
GO
SAVE 1 \$2.COM



a subsequent SUBMIT command, such as

```
A>SUBMIT SAVER PIP Y
```

substitutes X for \$1 and Y for \$2 in the command stream. The XSUB program loads, followed by DDT, which is sent to the command lines PIP.COM, R, and G0, thus returning to the CCP. The final command SAVE 1 Y.COM is processed by the CCP.

The XSUB program remains in memory and prints the message

```
(xsub active)
```

on each warm start operation to indicate its presence. Subsequent SUBMIT command streams do not require the XSUB, unless an intervening cold start occurs. Note that XSUB must be loaded after the optional CP/M DESPOOL utility, if both are to run simultaneously.

#### 1.6.8 DUMP Command

Syntax:

```
DUMP ufn
```

The DUMP program types the contents of the disk file (ufn) at sixteen bytes at a time, with the absolute byte address listed to the left of each line in hexadecimal. Long typeouts can be aborted by pressing the rubout key during printout. The source listing of the DUMP program is given in Section 5 as an example of a program written for the CP/M environment.

#### 1.6.9 MOVCPM Command

Syntax:

```
MOVCPM
```

The MOVCPM program allows you to reconfigure the CP/M system for any particular memory size. Two optional parameters can be used to indicate the desired size of the new system and the disposition of the new system at program termination. If the first parameter is omitted or an \* is given, the MOVCPM program reconfigures the system to its maximum size, based upon the kilobytes of contiguous RAM in the host system (starting at 0000H). If the second parameter is omitted, the system is executed, but not permanently recorded; if \* is given, the system is left in memory, ready for a SYSGEN operation. The MOVCPM program relocates a memory image of CP/M and places this image in memory in preparation for a system generation operation. The following is a list of MOVCPM command forms:



## Transient Commands







B	A	D		S	E	C	T	O	R
S	E	L	E	C	T				
R	E	A	D		O	N	L	Y	

You can also encounter this condition in reading files generated by a controller produced by a different manufacturer. Even though controllers claim to be IBM--compatible, one often finds small differences in recording formats. The Model 800 controller, for example, requires two bytes of one's following the data CRC byte, which is not required in the IBM format. As a result, disks generated by the Intel microcomputer development system can be read by almost all other IBM-compatible system, while disk files generated on other manufacturers' equipment produce the BAD SECTOR message when read by the microcomputer development system. To recover from this condition, press a CTRL-C to reboot (the safest course), or a return, which ignores the bad sector in

```

The READ ONLY message occurs when there is an attempt to write
to a disk or file that has been designated as Read-Only in a STAT
command or has been set to Read-Only by the BDOS.  Reboot CP/M by
using the warm start procedure, CTRL-C, or by performing a cold
start whenever the disks are changed.  If a changed disk is to be
read but not written, BDOS allows the disk to be changed without the
warm or cold start, but internally marks the drive as Read-Only.
The status of the drive is subsequently changed to Read-Write if a
warm or cold start occurs.  On issuing this message, CP/M waits for
input from the console.  An automatic warm start takes place
following any input.

```

1	-	3	8
---	---	---	---



This section gives operating procedures for using CP/M on the Intel Model 800 microcomputer development system microcomputer development system. Basic knowledge of the microcomputer development system hardware and software systems is assumed.

CP/M is initiated in essentially the same manner as the Intel ISIS operating system. The disk drives are labeled 0 through 3 on the microcomputer development system, corresponding to CP/M drives A through D, respectively. The CP/M system disk is inserted into drive 0, and the BOOT and RESET switches are pressed in sequence. The interrupt 2 light should go on at this point. The space bar is then pressed on the system console, and the light should go out. If it does not, the user should check connections and baud rates. The BOOT switch is turned off, and the CP/M sign-on message should appear at the selected console device, followed by the A> system prompt. You can then issue the various resident and transient commands.

The CP/M system can be restarted (warm start) at any time by pushing the INT 0 switch on the front panel. The built-in Intel ROM monitor can be initiated by pushing the INT 7 switch, which generates an RST 7, except when operating under DDT, in which case the DDT program gets control instead.

Diskettes can be removed from the drives at any time, and the system can be shut down during operation without affecting data integrity. Do not remove a disk and replace it with another without is Read-Only.

As a result of hardware hang-ups or malfunctions, CP/M might print the following message:

```
BDOS ERR ON d: BAD SECTOR
```

where d is the drive that has a permanent error. This error can occur when drive doors are opened and closed randomly, followed by disk operations, or can be caused by a disk, drive, or controller failure. You can optionally elect to ignore the error by pressing a single return at the console. The error might produce a bad data record, requiring reinitialization of up to 128 bytes of data. You can reboot the CP/M system and try the operation again.

Termination of a CP/M session requires no special action, except that it is necessary to remove the disks before turning the power off to avoid random transients that often make their way to the drive electronics.

You should use IBM-compatible disks rather than disks that have previously been used with any ISIS version. In particular, the ISIS FORMAT operation produces nonstandard sector numbering throughout the disk. This nonstandard numbering seriously degrades the performance of CP/M, and causes CP/M to operate noticeably slower than the distribution version. If it becomes necessary to reformat a disk, which should not be the case for standard disks, a program



Generally,	IBM-compatible	8-inch	disks	do	not	need	to	be
------------	----------------	--------	-------	----	-----	------	----	----

End of Section 1

	LINK MECHANISM	

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--	--	--

--	--	--

[illegible]

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--	--	--

--	--	--

--	--	--

--	--	--

--	--	--

[illegible][illegible]



## Section 2

### The CP/M Editor

#### 2.1 Introduction to Ed

Ed is the context editor for CP/M, and is used to create and alter CP/M source files. To start ED, type a command of the following form:

ED filename

or

ED filename.typ

Generally, ED reads segments of the source file given by filename or filename.typ into the central memory, where you edit the file and it is subsequently written back to disk after alterations. If the source file does not exist before editing, it is created by ED and initialized to empty. The overall operation of Ed is shown in Figure 2-1.

##### 2.1.1 ED Operation

Ed operates upon the source file, shown in Figure 2-1 by x.y, viewed or altered. The number of lines that can be maintained in the memory buffer varies with the line length, but has a total capacity of about 5000 characters in a 20K CP/M system.

Edited text material is written into a temporary work file under your command. Upon termination of the edit, the memory buffer is written to the temporary file, followed by any remaining (unread) text in the source file. The name of the original file is changed from x.y to x.BAK so that the most recent edited source file can be reclaimed if necessary. See the CP/M commands ERASE and RENAME. The temporary file is then changed from x.\$\$\$ to x.y, which becomes the resulting edited file.

The memory buffer is logically between the source file and working file, as shown in Figure 2-2.







## 2.1.2 Text Transfer Functions

Given that *n* is an integer value in the range 0 through 65535, several single-letter ED commands transfer lines of text from the source file through the memory buffer to the temporary (and eventually final) file. Single letter commands are shown in upper-case, but can be typed in either upper- or lower-case.

Table 2-1. ED Text Transfer Commands

Command	Result
<i>nA</i>	Appends the next <i>n</i> unprocessed source lines from the source file at SP to the end of the memory buffer at MP. Increment SP and MP by <i>n</i> . If upper-case translation is set (see the U command) and the A command is typed in upper-case, all input lines will automatically be translated to upper-case.
<i>nW</i>	Writes the first <i>n</i> lines of the memory buffer to the temporary file free space. Shift the remaining lines <i>n</i> +1 through MP to the top of the memory buffer. Increment TP by <i>n</i> .
E	Ends the edit. Copy all buffered text to temporary file. Rename files.
H	Moves to head of new file by performing automatic E command. The temporary file becomes the new source file, the memory buffer is emptied, and a new temporary file is created. The effect is equivalent to issuing an E command, followed by a reinvocation of ED, using <i>x.y</i> as the file to edit.
O	Returns to original file. The memory buffer is emptied, the temporary file is deleted, and the SP is returned to position 1 of the source file. The effects of the previous editing commands are thus nullified.
Q	Quits edit with no file alterations, returns to CP/M.

There are a number of special cases to consider. If the integer *n* is omitted in any ED command where an integer is allowed, then 1 is assumed. Thus, the commands A and W append one line and write one line, respectively. In addition, if a pound sign # is given in the place of *n*, then the integer 65535 is assumed (the largest value for *n* that is allowed). Because most source files can be contained entirely in the memory buffer, the command #A is often



issued at the beginning of the edit to read the entire source file to memory. Similarly, the command #W writes the entire buffer to the temporary file.

Two special forms of the A and W commands are provided as a convenience. The command 0A fills the current memory buffer at least half full, while 0W writes lines until the buffer is at least half empty. An error is issued if the memory buffer size is exceeded. You can then enter any command, such as W, that does not increase memory requirements. The remainder of any partial line read during the overflow will be brought into memory on the next successful append.

### 2.1.3 Memory Buffer Organization

The memory buffer can be considered a sequence of source lines brought in with the A command from a source file. The memory buffer has an imaginary character pointer (CP) that moves throughout the memory buffer under command of the operator.

The memory buffer appears logically as shown in Figure 2-3, where the dashes represent characters of the source line of indefinite length, terminated by carriage return (<cr>) and line-feed (<lf>) characters, and CP represents the imaginary character pointer. Note that the CP is always located ahead of the first character of the first line, behind the last character of the last line that contains the CP.

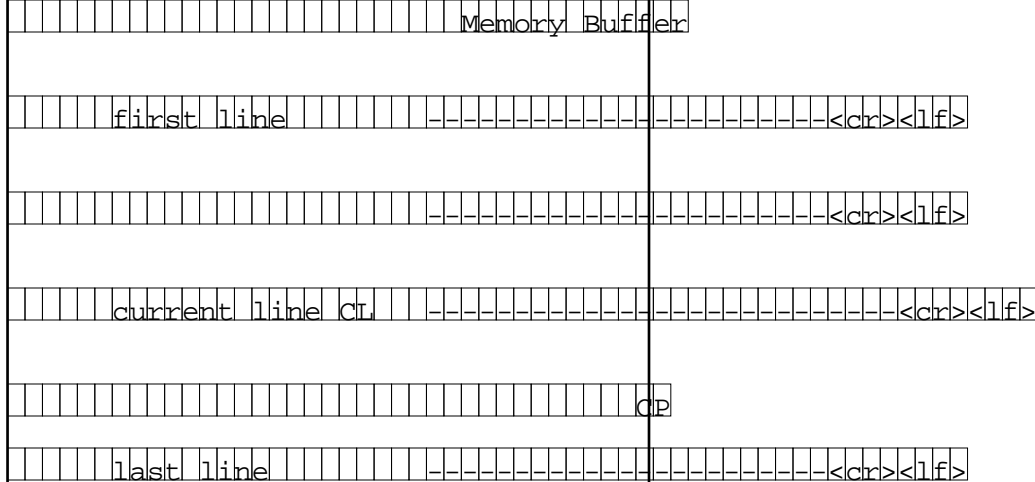


Figure 2-3. Logical Organization of Memory Buffer



[illegible]



[illegible]



Table 2-2. Editing Commands

Command	Action
$\pm$ B	Move CP to beginning of memory buffer if + and to bottom if -.
$\pm$ nC	Move CP by $\pm$ n characters (moving ahead if +), counting the <cr><lf> as two characters.
$\pm$ nD	Delete n characters ahead of CP if plus and behind CP if minus.
$\pm$ nK	Kill (remove) $\pm$ n lines of source text using CP as the current reference. If CP is not at the beginning of the current line when K is issued, the characters before CP remain if + is specified, while the characters after CP remain if - is given in the command.
$\pm$ nL	If n = 0, move CP to the beginning of the current line, if it is not already there. If n $\neq$ 0, first move the CP to the beginning of the current line and then move it to the beginning of the line that is n lines down (if +) or up (if -). The CP will stop at the top or bottom of the memory buffer if too large a value of n

$\pm$ nT	If n = 0, type the contents of the current line up to CP. If n = 1, type the contents of the current line from CP to the end of the line. If n>1, type the current line along with n - 1 lines that follow, if + is specified. Similarly, if n>1 and - is given, type the previous n lines up to the CP. Any key can be depressed to abort long type-outs.
$\pm$ n	Equivalent to $\pm$ nLT, which moves up or down and types a single line.

### 2.1.6 Command Strings

Any number of commands can be typed contiguously (up to the capacity of the console buffer) and are executed only after you press the <cr>. Table 2-3 summarizes the CP/M console line-editing commands used to control the input command line.



Table 2-3. Line-editing Controls	
Control	Function
^A	Move cursor to beginning of line
^B	Move cursor one character to the left
^C	Move cursor to end of line
^D	Delete character under cursor
^E	Move cursor one character to the right
^F	Move cursor to beginning of next line
^G	Move cursor to end of next line
^H	Delete character under cursor
^I	Move cursor to beginning of line
^J	Move cursor to end of line
^K	Delete character under cursor
^L	Move cursor to beginning of line
^M	Move cursor to end of line
^N	Move cursor to beginning of next line
^O	Move cursor to end of next line
^P	Move cursor to beginning of line
^Q	Move cursor to end of line
^R	Move cursor to beginning of line
^S	Move cursor to end of line
^T	Move cursor to beginning of line
^U	Move cursor to end of line
^V	Move cursor to beginning of line
^W	Move cursor to end of line
^X	Move cursor to beginning of line
^Y	Move cursor to end of line
^Z	Move cursor to beginning of line
^_	Move cursor to end of line

Command	Result
CTRL-C	Reboots the CP/M system when typed at the start of a line.
CTRL-E	Physical end of line: carriage is returned, but line is not sent until the carriage return key is depressed.
CTRL-H	Backspaces one character position.
CTRL-J	Terminates current input (line-feed).
CTRL-M	Terminates current input (carriage return).
CTRL-R	Retypes current command line: types a clean line character deletion with rubouts.
CTRL-U	Deletes the entire line typed at the console.
CTRL-X	Same as CTRL-U.
CTRL-Z	Ends input from the console (used in PIP and ED).

```

the console.

```

Suppose the memory buffer contains the characters shown in the previous section, with the CP following the last character of the buffer. In the following example, the command strings on the left produce the results shown to the right. Use lower-case command letters to avoid automatic translation of strings to upper-case.

Command String	Effect
B2T<cr>	Move to beginning of the buffer and type two lines:
	NOW IS THE
	TIME FOR
	The result in the memory buffer is
	NOW IS THE<cr><lf>
	TIME FOR<cr><lf>
	ALL GOOD MEN<cr><lf>







2-10



2	-	1	1
---	---	---	---



CP/M Operating System Manual	2.1 Introduction to ED
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ED also provides a single command that combines the F and I commands to perform simple string substitutions. The command takes the following form:	
--	--

								n	S		s	1	^	Z	s	2	<	c	r	>
--	--	--	--	--	--	--	--	---	---	--	---	---	---	---	---	---	---	---	---	---

ok
----

A diagram of a 16-bit register divided into seven fields. From left to right, the fields are: a 4-bit field, a 4-bit field labeled 'nS', a 1-bit field labeled 's', a 1-bit field labeled '1', a 4-bit field labeled '^Z', a 1-bit field labeled 's', and a 1-bit field labeled '2'. The fields are separated by vertical lines.

and has exactly the same effect as applying the following command string a total of  $n$  times:

[illegible]

or
----

						F	S	1	^Z-kDIs	2	^Z
--	--	--	--	--	--	---	---	---	---------	---	----

where  $k$  is the length of the string. ED searches the memory buffer starting at the current position of CP and successively substitutes the second string for the first string until the end of buffer, or until the substitution has been performed  $n$  times.

As a convenience, a command similar to F is provided by FD that automatically appends and writes lines as the search proceeds. The form is

					n	N	s	<	Cr>										
T	2	2	4	F	G	7	8	D	A	I	2	3	4	F	G	7	8		

11	2
0	7

						n	N	S	^Z
--	--	--	--	--	--	---	---	---	----

which searches the entire source file for the nth occurrence of the strings (you should recall that F fails if the string cannot be found in the current buffer). The operation of the N command is precisely the same as F except in the case that the string cannot be found within the current memory buffer. In this case, the entire memory content is written (that is, an automatic #W is issued). Input lines are then read until the buffer is at least half full, or the entire source file is exhausted. The search continues in this manner until the string has been found n times, or until the source file has been completely transferred to the temporary file.

A final line editing function, called the juxtaposition command, takes the form

							n	J	s	1	^Z	s	2	^Z	s	3	<cr>
--	--	--	--	--	--	--	---	---	---	---	----	---	---	----	---	---	------

ok
----

						n	J	s	$\wedge$ Z	s	$\wedge$ Z	s	$\wedge$ Z
								1		2		3	

with the following action applied n times to the memory buffer: search from the current CP for the next occurrence of the string s<sub>1</sub>. If found, insert the string s<sub>2</sub>, and move CP to follow s<sub>2</sub>. Then delete all characters following CP up to, but not including, the string s<sub>3</sub>, leaving CP directly after s<sub>2</sub>. If s<sub>3</sub> cannot be found, then no deletion is made. If the current line is



2	-	1	3
---	---	---	---



[illegible]

Note that upon normal completion of the ED program through Q or E, the temporary LTB file is removed. If ED is aborted with a CTRL-C, the LTB file will exist if lines have been transferred, but will generally be empty (a subsequent ED invocation will erase the temporary file).

						n	M	C	S	<cr>
--	--	--	--	--	--	---	---	---	---	------

	MFGAMMA^Z-5DIDELTA^Z0TT<cr>
--	-----------------------------

MSGAMMA^ZDELTA^Z0TT&lt;cr&gt;

2	-	1	4
---	---	---	---



Table 2-4. Error Message Symbols

Symbol	Meaning
?	Unrecognized command.
>	Memory buffer full (use one of the commands D, K, N, S, or W to remove characters); F, N, or S strings too long.
#	Cannot apply command the number of times specified (for example, in F command).
O	Cannot open IIB file in R command.

If there is a disk error, CP/M displays the following message:

```
BDOS ERR on d: BAD SECTOR
```

You can choose to ignore the error by pressing RETURN at the console (in this case, the memory buffer data should be examined to see if they were incorrectly read), or you can reset the system with a CTRL-C and reclaim the back-up file if it exists. The file can be reclaimed by first typing the contents of the BAK file to ensure that it contains the proper information. For example, type the following:

```
TYPE x.BAK
```

where x is the file being edited. Then remove the primary file

```
ERA x.y
```

and rename the BAK file

```
REN x.y=x.BAK
```

The file can then be reedited, starting with the previous version.

ED also takes file attributes into account. If you attempt to edit a Read-Only file, the message

```
** FILE IS READ/ONLY **
```

appears at the console. The file can be loaded and examined, but cannot be altered. You must end the edit session and use STAT to change the file attribute to R/W. If the edited file has the system attribute set, the following message:

```
'SYSTEM' FILE NOT ACCESSIBLE
```

is displayed and the edit session is aborted. Again, the STAT program can be used to change the system attribute, if desired.







32	

[illegible]







## Section 3

### CP/M Assembler

#### 3.1 Introduction

The CP/M assembler reads assembly-language source files from the disk and produces 8080 machine language in Intel hex format. To start the CP/M assembler, type a command in one of the following forms:

```
ASM filename
ASM filename.parms
```

In both cases, the assembler assumes there is a file on the disk with the name:

```
filename.ASM
```

which contains an 8080 assembly-language source file. The first and second forms shown above differ only in that the second form allows parameters to be passed to the assembler to control source file access and hex and print file destinations.

In either case, the CP/M assembler loads and prints the

```
CP/M ASSEMBLER VER n.n
```

where n.n is the current version number. In the case of the first command, the assembler reads the source file with assumed filetype ASM and creates two output files

```
filename.HEX
filename.PRN
```

The HEX file contains the machine code corresponding to the original program in Intel hex format, and the PRN file contains an annotated listing showing generated machine code, error flags, and source lines. If errors occur during translation, they are listed in the PRN file and at the console.

The form ASM filename parms is used to redirect input and output files from their defaults. In this case, the parms portion of the command is a three-letter group that specifies the origin of the source file, the destination of the hex file, and the destination of the print file. The form is



where  $p_1$ ,  $p_2$ , and  $p_3$  are single letters.  $P_1$  can be

								X
--	--	--	--	--	--	--	--	---

3	-	2
---	---	---



### 3.1 Introduction

### 3.2 Program Format

line#	label	operation	operand	;comment
-------	-------	-----------	---------	----------

The label field takes either of the following forms:

The label field is optional, except where noted in particular statement types. The identifier is a sequence of alphanumeric characters where the first character is alphabetic. Identifiers can be freely used by the programmer to label elements such as program steps and assembler directives, but cannot exceed 16 characters in length. All characters are significant in an identifier, except for the embedded dollar symbol \$, which can be used to improve readability of the name. Further, all lower-case alphabetic characters are treated as upper-case. The following are all valid instances of labels:

x1	y2	x1	x2	x2	34	\$5	67	8\$	90	12	\$3	45	6:
----	----	----	----	----	----	-----	----	-----	----	----	-----	----	----

3-3



Generally, the operand field of the statement contains an expression formed out of constants and labels, along with arithmetic and logical operations on these elements. Again, the complete details of properly formed expressions are given in Section 3.3.

The comment field contains arbitrary characters following the semicolon symbol until the next real or logical end-of-line. These characters are read, listed, and otherwise ignored by the assembler. The CP/M assembler also treats statements that begin with an \* in column one as comment statements that are listed and ignored in the assembly process.

The assembly-language program is formulated as a sequence of statements of the above form, terminated by an optional END statement. All statements following the END are ignored by the assembler.

### 3.3 Forming the Operand

To describe the operation codes and pseudo operations completely, it is necessary first to present the form of the operand field, since it is used in nearly all statements. Expressions in the operand field consist of simple operands, labels, constants, and reserved words, combined in properly formed subexpressions by arithmetic and logical operators. The expression computation is carried out by the assembler as the assembly proceeds. Each

Further, the number of significant digits in the result must not exceed the intended use. If an expression is to be used in a byte move immediate instruction, the most significant 8 bits of the expression must be zero. The restriction on the expression significance is given with the individual instructions.

#### 3.3.1 Labels

As discussed above, a label is an identifier that occurs on a particular statement. In general, the label is given a value determined by the type of statement that it precedes. If the label occurs on a statement that generates machine code or reserves memory space (for example, a MOV instruction or a DS pseudo operation), the label is given the value of the program address that it labels. If the label precedes an EQU or SET, the label is given the value that results from evaluating the operand field. Except for the SET statement, an identifier can label only one statement.

When a label appears in the operand field, its value is substituted by the assembler. This value can then be combined with other operands and operators to form the operand field for a particular instruction.



CP/M Operating System Manual 3.3 Forming the Operand

### 3.3.2 Numeric Constants

A numeric constant is a 16-bit value in one of several bases. The base, called the radix of the constant, is denoted by a trailing radix indicator. The following are radix indicators:

- o B is a binary constant (base 2).
- o O is a octal constant (base 8).
- o Q is a octal constant (base 8).
- o D is a decimal constant (base 10).
- o H is a hexadecimal constant (base 16).

Q is an alternate radix indicator for octal numbers because the letter O is easily confused with the digit 0. Any numeric constant that does not terminate with a radix indicator is a decimal constant.

A constant is composed as a sequence of digits, followed by an optional radix indicator, where the digits are in the appropriate range for the radix. Binary constants must be composed of 0 and 1 digits, octal constants can contain digits in the range 0-7, while decimal constants contain decimal digits. Hexadecimal constants contain decimal digits as well as hexadecimal digits A(10D), B(11D), C(12D), D(13D), E(14D), and F(15D). Note that the leading digit of a hexadecimal constant must be a decimal digit to avoid confusing a hexadecimal constant with an identifier. A leading 0 will always suffice. A constant composed in this manner must evaluate to a

otherwise it is truncated on the right by the assembler.

Similar to identifiers, embedded \$ signs are allowed within constants to improve their readability. Finally, the radix indicator is translated to upper-case if a lower-case letter is encountered. The following are all valid instances of numeric constants:

1234	1234D	1100B	1111\$0000\$1111\$0000B
1234H	OFFEH	3377O	33\$77\$22Q
3377o	ofe3h	1234d	offffh

### 3.3.3 Reserved Words

There are several reserved character sequences that have predefined meanings in the operand field of a statement. The names of 8080 registers are given below. When they are encountered, they produce the values shown to the right.



Table 3-1. Reserved Characters

Character	Value
A	7
B	0
C	1
D	2
E	3
H	4
L	5
M	6
SP	6
PSW	6

Again, lower-case names have the same values as their upper-case equivalents. Machine instructions can also be used in the operand field; they evaluate to their internal codes. In the case of instructions that require operands, where the specific operand becomes a part of the binary bit pattern of the instruction, for example, MOV A,B, the value of the instruction, in this case MOV, is the bit pattern of the instruction with zeros in the optional fields, for example, MOV produces 40H.

When the symbol \$ occurs in the operand field, not embedded within identifiers and numeric constants, its value becomes the instruction contained within the current logical line.

### 3.3.4 String Constants

String constants represent sequences of ASCII characters and are represented by enclosing the characters within apostrophe symbols. All strings must be fully contained within the current physical line (thus allowing exclamation point symbols within strings) and must not exceed 64 characters in length. The apostrophe character itself can be included within a string by representing it as a double apostrophe (the two keystrokes ''), which becomes a single apostrophe when read by the assembler. In most cases, the string length is restricted to either one or two characters (the DB pseudo operation is an exception), in which case the string becomes an 8- or 16-bit value, respectively. Two-character strings become a 16-bit constant, with the second character as the low-order byte, and the first character as the high-order byte.

The value of a character is its corresponding ASCII code. There is no case translation within strings; both upper- and lower-case characters can be represented. You should note that only graphic printing ASCII characters are allowed within strings.



I said "Hello" to her.	I said "Hello" to her.
------------------------	------------------------

	3	-	7



### 3.3 Forming the Operand

[illegible]



(	a	*	b	)	+	c
---	---	---	---	---	---	---

a	+	(	b	*	c	)
---	---	---	---	---	---	---

(	(	a		M	O	D		b	)	*	c	)		S	H	L		d
---	---	---	--	---	---	---	--	---	---	---	---	---	--	---	---	---	--	---

```
a OR (b AND (NOT (c+(d SHL e))))
```

SHIL e

$$(d \text{ SHL } e)$$

### 3.4 Assembler Directives

Table 3-3. Assembler Directives																			
---------------------------------	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--

					M	e	a	n	i	n	g
--	--	--	--	--	---	---	---	---	---	---	---

ram or data origin

optional	start	address
----------	-------	---------

e	
---	--

ional assembly

ional assembly
----------------

out es

words

storage area



CP/M Operating System Guide 3.4 Assembler Directives

### 3.4.1 The ORG Directive

The ORG statement takes the form:

label ORG expression

where label is an optional program identifier and expression is a 16-bit expression, consisting of operands that are defined before the ORG statement. The assembler begins machine code generation at the location specified in the expression. There can be any number of ORG statements within a particular program, and there are no checks to ensure that the programmer is not defining overlapping memory areas. Note that most programs written for the CP/M system begin with an ORG statement of the form:

ORG 100H

which causes machine code generation to begin at the base of the CP/M transient program area. If a label is specified in the ORG statement, the label is given the value of the expression. This label can then be used in the operand field of other statements to represent this expression.

### 3.4.2 The END Directive

The END statement is optional in an assembly-language program. statements are ignored in the assembly. The END statement takes the following two forms:

label END

label END expression

where the label is again optional. If the first form is used, the assembly process stops, and the default starting address of the program is taken as 0000. Otherwise, the expression is evaluated, and becomes the program starting address. This starting address is included in the last record of the Intel-formatted machine code hex file that results from the assembly. Thus, most CP/M assembly-language programs end with the statement:

END 100H

resulting in the default starting address of 100H (beginning of the transient program area).



CP/M Operating System Guide 3.4 Assembler Directives

### 3.4.3 The EQU Directive

The EQU (equate) statement is used to set up synonyms for particular numeric values. The EQU statement takes the form:

label EQU expression

where the label must be present and must not label any other statement. The assembler evaluates the expression and assigns this value to the identifier given in the label field. The identifier is usually a name that describes the value in a more human-oriented manner. Further, this name is used throughout the program to place parameters on certain functions. Suppose data received from a teletype appears on a particular input port, and data is sent to the teletype through the next output port in sequence. For example, you can use this series of equate statements to define these ports for a particular hardware environment:

TTYBASE EQU 10H ;BASE PORT NUMBER FOR TTY

TTYIN EQU TTYBASE ;TTY DATA IN

TTYOUT EQU TTYBASE+1 ;TTY DATA OUT

At a later point in the program, the statements that access the

IN TTYIN ;READ TTY DATA TO REG-A

....

OUT TTYOUT ;WRITE DATA TO TTY FROM REG-A

making the program more readable than if the absolute I/O ports are used. Further, if the hardware environment is redefined to start the teletype communications ports at 7FH instead of 10H, the first statement need only be changed to

TTYBASE EQU 7FH ;BASE PORT NUMBER FOR TTY

and the program can be reassembled without changing any other statements.

### 3.4.4 The SET Directive

The SET statement is similar to the EQU, taking the form:

label SET expression

except that the label can occur on other SET statements within the program. The expression is evaluated and becomes the current value



\_\_\_\_\_

\_\_\_\_\_

\_\_\_\_\_

\_\_\_\_\_

\_\_\_\_\_

\_\_\_\_\_

\_\_\_\_\_

\_\_\_\_\_

6789012345678901234



```

CP/M Operating System Guide      3.4 Assembler Directives

;
; TRUE EQU OFFFHH ;DEFINE VALUE OF TRUE
; FALSE EQU NOT TRUE ;DEFINE VALUE OF FALSE
;
; TTY EQU TRUE ;TRUE IF TTY, FALSE IF CRT
;
; TTYBASE EQU 10H ;BASE OF TTY I/O PORTS
; CRTBASE EQU 20H ;BASE OF CRT I/O PORTS
; IF TTY ;ASSEMBLE RELATIVE TO
; ;TTTYBASE
; CONIN EQU TTYBASE ;CONSOLE INPUT
; CONOUT EQU TTYBASE+1 ;CONSOLE OUTPUT
; ENDIF

; IF NOT TTY ;ASSEMBLE RELATIVE TO
; ;CRTBASE
; CONIN EQU CRTBASE ;CONSOLE INPUT
; CONOUT EQU CRTBASE+1 ;CONSOLE OUTPUT
; ENDIF

;
; IN CONIN ;READ CONSOLE DATA
; OUT CONOUT ;WRITE CONSOLE DATA

```

In this case, the program assembles for an environment where a teletype is connected, based at port 10H. The statement defining

```

; TTY EQU FALSE

```

and, in this case, the program assembles for a CRT based at port 20H.

### 3.4.6 The DB Directive

The DB directive allows the programmer to define initialized storage areas in single-precision byte format. The DB statement takes the form:

```

label DB e#1, e#2, ..., e#n

```

where e#1 through e#n are either expressions that evaluate to 8-bit values (the high-order bit must be zero) or are ASCII strings of length no greater than 64 characters. There is no practical restriction on the number of expressions included on a single source line. The expressions are evaluated and placed sequentially into the machine code file following the last program address generated by the assembler. String characters are similarly placed into memory starting with the first character and ending with the last character. Strings of length greater than two characters cannot be used as operands in more complicated expressions.



**Note:** ASCII characters are always placed in memory with the parity bit reset (0). Also, there is no translation from lower- to upper-case within strings. The optional label can be used to reference the data area throughout the remainder of the program. The following are examples of valid DB statements:

```
data:      DB      0,1,2,3,4,5
           DB      data and 0ffh,5,377Q,1+2+3+4

sign-on:   DB      'please type your name',cr,lf,0
           DB      'AB' SHR 8, 'C', 'DE' AND 7FH
```

### 3.4.7 The DW Directive

The DW statement is similar to the DB statement except double-precision two-byte words of storage are initialized. The DW statement takes the form:

```
label      DW      e#1, e#2, ..., e#n
```

where e#1 through e#n are expressions that evaluate to 16-bit results. Note that ASCII strings of one or two characters are allowed, but strings longer than two characters are disallowed. In all cases, the data storage is consistent with the 8080 processor; memory, followed by the most significant byte. The following are examples of DW statements:

```
doub:      DW      0ffefh,doub+4,signon-$,255+255
           DW      'a', 5, 'ab', 'CD', 6 shl 8 or 11b.
```

### 3.4.8 The DS Directive

The DS statement is used to reserve an area of uninitialized memory, and takes the form:

```
label      DS      expression
```

where the label is optional. The assembler begins subsequent code generation after the area reserved by the DS. Thus, the DS statement given above has exactly the same effect as the following statement:

```
label:     EQU     $ ;LABEL VALUE IS CURRENT CODE LOCATION
           ORG     $+expression ;MOVE PAST RESERVED AREA
```



3.5 Operation Codes

Assembly-language operation codes form the principal part of assembly-language programs and form the operation field of the instruction. In general, ASM accepts all the standard mnemonics for the Intel 8080 microcomputer, which are given in detail in the Intel 8080 Assembly Language Programming Manual. Labels are optional on each input line. The individual operators are listed briefly in the following sections for completeness, although the Intel manuals should be referenced for exact operator details. In Tables 3-4 through 3-8, bit values have the following meaning:

o e3 represents a 3-bit value in the range 0-7 that can be one of the predefined registers A, B, C, D, E, H, L, M, SP, or PSW.

o e8 represents an 8-bit value in the range 0-255.

o e16 represents a 16-bit value in the range 0-65535.

These expressions can be formed from an arbitrary combination of operands and operators. In some cases, the operands are restricted to particular values within the allowable range, such as the PUSH instruction. These cases are noted as they are encountered.

its most general form, along with a specific example, a short explanation, and special restrictions.

3.5.1 Jumps, Calls, and Returns

The Jump, Call, and Return instructions allow several different forms that test the condition flags set in the 8080 microcomputer CPU. The forms are shown in Table 3-4.

Table 3-4. Jumps, Calls, and Returns

Form	Bit Value	Example	Meaning
JMP	e16	JMP L1	Jump unconditionally to label
JNZ	e16	JNZ L2	Jump on nonzero condition to label
JZ	e16	JZ 100H	Jump on zero condition to label
JNC	e16	JNC L1+4	Jump no carry to label
JC	e16	JC L3	Jump on carry to label
JPO	e16	JPO \$+8	Jump on parity odd to label



Table 3-4. (continued)

Form	Bit	Example	Meaning
	Value		
JPE	e16	JPE L4	Jump on even parity to label
JP	e16	JP GAMMA	Jump on positive result to label
JM	e16	JM al	Jump on minus to label
CALL	e16	CALL S1	Call subroutine unconditionally
CNZ	e16	CNZ S2	Call subroutine on nonzero condition
CZ	e16	CZ 100H	Call subroutine on zero condition
CNC	e16	CNC S1+4	Call subroutine if no carry set
CC	e16	CC S3	Call subroutine if carry set
CPO	e16	CPO \$+8	Call subroutine if parity odd
CPE	e16	CPE \$4	Call subroutine if parity even

CM	e16	CM b1\$c2	Call subroutine if minus flag
RST	e3	RST 0	Programmed restart, equivalent to CALL 8*e3, except one byte call
RET			Return from subroutine
RNZ			Return if nonzero flag set
RZ			Return if zero flag set
RNC			Return if no carry
RC			Return if carry flag set
RPO			Return if parity is odd
RPE			Return if parity is even
RP			Return if positive result
RM			Return if minus flag is set



3.5.2 Immediate Operand Instructions

Several instructions are available that load single- or double-precision registers or single-precision memory cells with constant values, along with instructions that perform immediate arithmetic or logical operations on the accumulator (register A). Table 3-5 describes the immediate operand instructions.

Table 3-5. Immediate Operand Instructions

Form with Bit Values	Example	Meaning
MVI e3,e8	MVI B,255	Move immediate data to register A, B, C, D, E, H, L, or M (memory)
ADI e8	ADI 1	Add immediate operand to A without carry
ACI e8	ACI 0FFH	Add immediate operand to A with carry
SUI e8	SUI L + 3	Subtract from A without borrow (carry)
		(carry)
ANI e8	ANI \$ AND 7FH	Logical and A with immediate data
XRI e8	XRI 1111\$0000B	Exclusive or A with immediate data
ORI e8	ORI L AND 1+1	Logical or A with immediate data
CPI e8	CPI 'a'	Compare A with immediate data, same as SUI except register A not changed.
LXI e3,e16	LXI B,100H	Load extended immediate to register pair. e3 must be equivalent to B, D, H, or SP.

3.5.3 Increment and Decrement Instructions

The 8080 provides instructions for incrementing or decrementing single- and double-precision registers. The instructions are described in Table 3-6.



Table 3-6. Increment and Decrement Instructions

Form with Bit Value	Example	Meaning
INR e3	INR E	Single-precision increment register. e3 produces one of A, B, C, D, E, H, L, M.
DCR e3	DCR A	Single-precision decrement register. e3 produces one of A, B, C, D, E, H, L, M.
INX e3	INX SP	Double-precision increment register pair. e3 must be equivalent to B, D, H, or SP.
DCX e3	DCX B	Double-precision decrement register pair. e3 must be equivalent to B, D, H, or SP.

3.5.4 Data Movement Instructions

to memory are given in the following table.

Table 3-7. Data Movement Instructions

Form with Bit Value	Example	Meaning
MOV e3,e3	MOV A,B	Move data to leftmost element from rightmost element. e3 produces one of A, B, C, D, E, H, L, or M. MOV M,M is disallowed.
LDAX e3	LDAX B	Load register A from computed address. e3 must produce either B or D.
STAX e3	STAX D	Store register A to computed address. e3 must produce either B or D.
LHLD e16	LHLD L1	Load HL direct from location e16. Double-precision load to H and L.



Meaning
---------

Load register A with data

Exchange	DE	pair	with	HL
pair				

## 3-19







Table 3-8. (continued)

Form with Bit Value	Example	Meaning
RAL		Rotate carry/A register to left. Carry is involved in the rotate.
RAR		Rotate carry/A register to right. Carry is involved in the rotate.
DAD e3	DAD B	Double-precision add register pair e3 to HL. e3 must produce B, D, H, or SP.

3.5.6 Control Instructions

The four remaining instructions, categorized as control instructions, are the following:

- o HLT halts the 8080 processor.
- o DI disables the interrupt system.
- o EI enables the interrupt system.

3.6 Error Messages

When errors occur within the assembly-language program, they are listed as single-character flags in the leftmost position of the source listing. The line in error is also echoed at the console so that the source listing need not be examined to determine if errors are present. The error codes are listed in the following table.

Table 3-9. Error Codes

Error Code	Meaning
D	Data error: element in data statement cannot be placed in the specified data area.
E	Expression error: expression is ill-formed and cannot be computed at assembly time.
L	Label error: label cannot appear in this context; might be duplicate label.
N	Not implemented: features that will appear in future ASM versions. For example, macros are recognized, but flagged in this version.







3	-	2	3
---	---	---	---



[illegible]



[illegible]



	CP/M Operating System Manual	3.7 A Sample Session
C0Z0M1E0I0 A=00 B=0005 D=0000 H=0148 S=0100 P=0122 MOV A,C		
C0Z0M1E0I0 A=05 B=0005 D=0000 H=0148 S=0100 P=0123 INX H		
C0Z0M1E0I0 A=05 B=0005 D=0000 H=0149 S=0100 P=0124 MOV B,M*0125	-L100 Automatic breakpoint	
IXT H,0146	MVI M,01	LXI H,0147
MVI M,00	MOV A,M List some code	CPI 09 from 100H
JC 0119	LXI H,0146	MOV A,M
ORA A	JNZ 0110	-T-
RST 07	MOV E,A List more	MVI D,00
LXI H,0148	-Abort list with rubout	-G,11B Start program from current PC (0125H)
		and run in real time to 11BH
*0127 Stopped with an external interrupt 7 from front panel	-T4 (program was looping indefinitely)	Look at looping program in trace mode
C0Z0M0E0I0 A=38 B=0064 D=0006 H=0156 S=0100 P=0127 MOV D,A		
C0Z0M0E0I0 A=38 B=0064 D=3806 H=0156 S=0100 P=0128 MOV A,B		
C0Z0M0E0I0 A=00 B=0064 D=3806 H=0156 S=0100 P=0129 INX H		
C0Z0M0E0I0 A=00 B=0064 D=3806 H=0157 S=0100 P=012A SBB M*012B	-D148 Data are sorted, but program does not stop.	
	0148 05 00 07 00 14 00 1E 00.....	0150 32 00 64 00 64 00 2C 01 E8 03 01 80 00 00 00 00 2.D.D.,.....
	0160 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 .....	
-G0 Return to CP/M	A>DDT SORT.HEX Reload the memory image	
16K DDT VER 1.0	NEXT PC	015C 0000
-XP	P=0000 100 Set PC to beginning of program	-L10D List bad OPCODE
		3-26



[illegible]



[illegible]



[illegible]

## CP/M Dynamic Debugging Tool

## 4.1 Introduction

The DDT program allows dynamic interactive testing and debugging of programs generated in the CP/M environment. Invoke the debugger with a command of one of the following forms:

[illegible]

where filename is the name of the program to be loaded and tested. In both cases, the DDT program is brought into main memory in place of the Console Command Processor (CCP) and resides directly below the Basic Disk Operating System (BDOS) portion of CP/M. Refer to Section 5 for standard memory organization. The BDOS starting address, located in the address field of the JMP instruction at location 5H, is altered to reflect the reduced Transient Program Area (TPA) size.

The second and third forms of the DDT command perform the same actions as the first, except there is a subsequent automatic load of

following sequence of commands:

[illegible]

where the I and R commands set up and read the specified program to test. See the explanation of the I and R commands below for exact details.

Upon initiation, DDT prints a sign-on message in the form:

[illegible]

where  $m.m$  is the revision number.

Following the sign-on message, DDT prompts you with the hyphen character, -, and waits for input commands from the console. You can type any of several single-character commands, followed by a carriage return to execute the command. Each line of input can be line-edited using the following standard CP/M controls:



Table 4-1. Line-editing Controls

Control	Result
rubout	removes the last character typed
CTRL-U	removes the entire line, ready for retyping
CTRL-C	reboots system

Any command can be up to 32 characters in length. An automatic carriage return is inserted as character 33, where the first character determines the command type. Table 4-2 describes DDT commands.

Table 4-2. DDT Commands

Command Character	Result
A	enters assembly-language mnemonics with operands.
D	displays memory in hexadecimal and ASCII.
G	begins execution with optional breakpoints.
I	sets up a standard input File Control Block.
L	lists memory using assembler mnemonics.
M	moves a memory segment from source to destination.
R	reads a program for subsequent testing.
S	substitutes memory values.
T	traces program execution.
U	untraced program monitoring.
X	examines and optionally alters the CPU state.

The command character, in some cases, is followed by zero, one, two, or three hexadecimal values, which are separated by commas or single blank characters. All DDT numeric output is in hexadecimal form. The commands are not execution until the carriage return is typed at the end of the command.



At any point in the debug run, you can stop execution of DDT by using either a CTRL-C or G0 (jump to location 0000H) and save the current memory image by using a SAVE command of the form:

```
SAVE n filename.COM
```

where n is the number of pages (256 byte blocks) to be saved on disk. The number of blocks is determined by taking the high-order byte of the address in the TPA and converting this number to decimal. For example, if the highest address in the TPA is 134H, the number of pages is 12H or 18 in decimal. You could type a CTRL-C during the debug run, returning to the CCP level, followed by

```
SAVE 18 X.COM
```

The memory image is saved as X.COM on the disk and can be directly executed by typing the name X. If further testing is required, the memory image can be recalled by typing

```
DDT X.COM
```

which reloads the previously saved program from location 100H through page 18, 23FFH. The CPU state is not a part of the COM file; thus, the program must be restarted from the beginning to test it properly.

#### DDT Commands

The individual commands are detailed below. In each case, the operator must wait for the hyphen prompt character before entering the command. If control is passed to a program under test, and the program has not reached a breakpoint, control can be returned to DDT by executing a RST 7 from the front panel. In the explanation of each command, the command letter is shown in some cases with numbers separated by commas, the the numbers are represented by lower-case letters. These numbers are always assumed to be in a hexadecimal radix and from one to four digits in length. Longer numbers are automatically truncated on the right.

Many of the commands operate upon a CPU state that corresponds to the program under test. The CPU state holds the registers of the program being debugged and initially contains zeros for all registers and flags except for the program counter, P, and stack pointer, S, which default to 100H. The program counter is subsequently set to the starting address given in the last record of a HEX file if a file of this form is loaded, see the I and R commands.

##### 4.2.1 The A (Assembly) Command

DDT allows in-line assembly language to be inserted into the current memory image using the A command, which takes the form:



As

where s is the hexadecimal starting address for the in-line assembly. DDT prompts the console with the address of the next instruction to fill and reads the console, looking for assembly-language mnemonics followed by register references and operands in absolute hexadecimal form. See the Intel 8080 Assembly Language Reference Card for a list of mnemonics. Each successive load address is printed before reading the console. The A command terminates when the first empty line is input from the console.

Upon completion of assembly language input, you can review the memory segment using the DDT disassembler (see the I command).

Note that the assembler/disassembler portion of DDT can be overlaid by the transient program being tested, in which case the DDT program responds with an error condition when the A and I commands are used.

#### 4.2.2 The D (Display) Command

The D command allows you to view the contents of memory in hexadecimal and ASCII formats. The D command takes the forms:

D

Ds

In the first form, memory is displayed from the current display address, initially 100H, and continues for 16 display lines. Each display line takes the following form:

aaaa bb bb bb bb bb bb bb bb bb bb bb bb bb bb cccccccccccccccc

where aaaa is the display address in hexadecimal and bb represents data present in memory starting at aaaa. The ASCII characters starting at aaaa are to the right (represented by the sequence of character c) where nongraphic characters are printed as a period. You should note that both upper- and lower-case alphabetics are displayed, and will appear as upper-case symbols on a console device that supports only upper-case. Each display line gives the values of 16 bytes of data, with the first line truncated so that the next line begins at an address that is a multiple of 16.

The second form of the D command is similar to the first, except that the display address is first set to address s.

The third form causes the display to continue from address s through address f. In all cases, the display address is set to the first address not displayed in this command, so that a continuing display can be accomplished by issuing successive D commands with no explicit addresses.



4-5



Execution continues from the starting address in real-time to the next breakpoint. There is no intervention between the starting address and the break address by DDT. If the program under test does not reach a breakpoint, control cannot return to DDT without executing a RST 7 instruction. Upon encountering a breakpoint, DDT stops execution and types

```
*d
```

where d is the stop address. The machine state can be examined at this point using the X (Examine) command. You must specify breakpoints that differ from the program counter address at the beginning of the G command. Thus, if the current program counter is 1234H, then the following commands:

```
G,1234
```

```
G400,400
```

both produce an immediate breakpoint without executing any instructions.

#### 4.2.5 The I (Input) Command

The I command allows you to insert a filename into the default File Control Block (FCB) at 5CH. The FCB created by CP/M for transient programs is placed at this location (see Section 5). The filename is passed by the CP/M Console Processor. Note that this filename is also used by DDT for reading additional HEX and COM files. The I command takes the forms:

```
Ifilename
```

```
Ifilename.typ
```

If the second form is used and the filetype is either HEX or COM, subsequent R commands can be used to read the pure binary or hex format machine code. Section 4.2.8 gives further details.

#### 4.2.6 The L (List) Command

The L command is used to list assembly-language mnemonics in a particular program region. The L command takes the forms:

```
L
```

```
Ls
```

```
Ls,f
```

The first form lists twelve lines of disassembled machine code from the current list address. The second form sets the list address to s and then lists twelve lines of code. The last form lists disassembled code from s through address f. In all three cases, the list address is set to the next unlisted location in preparation for a subsequent L command. Upon encountering an



#### 4.2.7 The M (Move) Command

					M	S,	f,	d
--	--	--	--	--	---	----	----	---

#### 4.2.8 The R (Read) Command

						R
						RB

```
DDT filename.filetype
```

[illegible]



Whenever the R command is issued, DDT responds with either the error indicator ? (file cannot be opened, or a checksum error occurred in a HEX file) or with a load message. The load message takes the form:

```
      NEXT PC
      nnnn pppp
```

where nnnn is the next address following the loaded program and pppp is the assumed program counter (100H for COM files, or taken from the last record if a HEX file is specified).

#### 4.2.9 The S (Set) Command

The S command allows memory locations to be examined and optionally altered. The S command takes the form:

```
      Ss
```

where s is the hexadecimal starting address for examination and alteration of memory. DDT responds with a numeric prompt, giving the memory location, along with the data currently held in memory. If you type a carriage return, the data is not altered. If a byte value is typed, the value is stored at the prompted address. In either case, DDT continues to prompt with successive addresses and detected.

#### 4.2.10 The T (Trace) Command

The T command allows selective tracing of program execution for 1 to 65535 program steps. The T command takes the forms:

```
      T
      Tn
```

In the first form, the CPU state is displayed and the next program step is executed. The program terminates immediately, with the termination address displayed as

```
      *hhhh
```

where hhhh is the next address to execute. The display address (used in the D command) is set to the value of H and I, and the list address (used in the I command) is set to hhhh. The CPU state at program termination can then be examined using the X command.

The second form of the T command is similar to the first, except that execution is traced for n steps (n is a hexadecimal value) before a program breakpoint occurs. A breakpoint can be forced in the trace mode by typing a rubout character. The CPU state is displayed before each program step is taken in trace mode. The format of the display is the same as described in the X command.



You should note that program tracing is discontinued at the CP/M interface and resumes after return from CP/M to the program under test. Thus, CP/M functions that access I/O devices, such as the disk drive, run in real-time, avoiding I/O timing problems. Programs running in trace mode execute approximately 500 times slower than real-time because DDT gets control after each user instruction is executed. Interrupt processing routines can be traced, but commands that use the breakpoint facility (G, T, and U) accomplish the break using an RST 7 instruction, which means that the tested program cannot use this interrupt location. Further, the trace mode always runs the tested program with interrupts enabled, which may cause problems if asynchronous interrupts are received during tracing.

To get control back to DDT during trace, press RETURN rather than executing an RST 7. This ensures that the trace for current instruction is completed before interruption.

4.2.11 The U (Untrace) Command

The U command is identical to the T command, except that intermediate program steps are not displayed. The untrace mode allows from 1 to 65535, (0FFFFH) steps to be executed in monitored mode and is used principally to retain control of an executing program while it reaches steady state conditions. All conditions of the T command apply to the U command.

4.2.12 The X (Examine) Command

The X command allows selective display and alteration of the current CPU state for the program under test. The X command takes the forms:

X  
Xr

where r is one of the 8080 CPU registers listed in the following table.

Table 4-3. CPU Registers

Register	Meaning	Value
C	Carry flag	(0/1)
Z	Zero flag	(0/1)
M	Minus flag	(0/1)
E	Even parity flag	(0/1)
I	Interdigit carry	(0/1)
A	Accumulator	(0-FF)
B	BC register pair	(0-FFFF)
D	DE register pair	(0-FFFF)



Table 4-3. (continued)

Register	Meaning	Value
H	HL register pair	(0-FFFF)
S	Stack pointer	(0-FFFF)
P	Program counter	(0-FFFF)

In the first case, the CPU register state is displayed in the format:

```
CfZfMfEfIf A=bb B=dddd D=dddd H=dddd S=dddd P=dddd inst
```

where f is a 0 or 1 flag value, bb is a byte value, and dddd is a double-byte quantity corresponding to the register pair. The inst field contains the disassembled instruction, that occurs at the location addressed by the CPU state's program counter.

The second form allows display and optional alteration of register values, where r is one of the registers given above (C, Z, M, E, I, A, B, D, H, S, or P). In each case, the flag or register value is first displayed at the console. The DDT program then accepts input from the console. If a carriage return is typed, the flag or register value is not altered. If a value in the proper range is typed, the flag or register value is altered. You should note that BC, DE, and HL are displayed as register pairs. Thus, you

altered.

### 4.3 Implementation Notes

The organization of DDT allows certain nonessential portions to be overlaid to gain a larger transient program area for debugging large programs. The DDT program consists of two parts: the DDT nucleus and the assembler/disassembler module. The DDT nucleus is loaded over the CCP and, although loaded with the DDT nucleus, the assembler/disassembler is overlayable unless used to assemble or disassemble.

In particular, the BDOS address at location 6H (address field of the JMP instruction at location 5H) is modified by DDT to address the base location of the DDT nucleus, which, in turn, contains a JMP instruction to the BDOS. Thus, programs that use this address field to size memory see the logical end of memory at the base of the DDT nucleus rather than the base of the BDOS.

The assembler/disassembler module resides directly below the DDT nucleus in the transient program area. If the A, I, T, or X commands are used during the debugging process, the DDT program again alters the address field at 6H to include this module, further reducing the logical end of memory. If a program loads beyond the beginning of the assembler/disassembler module, the A and I commands are lost (their use produces a ? in response) and the trace and



4	-	1	1
---	---	---	---



*BOP			
------	--	--	--

[illegible][illegible]

VECT:	DB	2,0,4,3,5,6,1,5	
LEN	EQU	\$-VECT	;LENGTH
LARGE:	DS	1	;LARGEST VALUE ON EXIT
	END		
*E <--End of edit			

```
A>ASM SCAN Start Assembler
```

CP/M ASSEMBLER - VER 1.0

0	1	2	2
---	---	---	---

0.02H	USE	FACTOR
-------	-----	--------

```
END OF ASSEMBLY      Assembly complete; look at program listing
```

```
A>TYPE SCAN.PRN
```

Code address	Source program
--------------	----------------

[illegible]



[illegible]



1234567890123456789012345678901234567890123456789012345678901







# A Sample Program

4-16



[illegible]



[illegible]



[illegible]



[illegible]



[illegible]



## Section 5

### CP/M 2 System Interface

#### 5.1 Introduction

This chapter describes CP/M (release 2) system organization including the structure of memory and system entry points. This section provides the information you need to write programs that operate under CP/M and that use the peripheral and disk I/O facilities of the system.

CP/M is logically divided into four parts, called the Basic Input/Output System (BIOS), the Basic Disk Operating System (BDOS), the Console Command Processor (CCP), and the Transient Program Area (TPA). The BIOS is a hardware-dependent module that defines the exact low level interface with a particular computer system that is necessary for peripheral device I/O. Although a standard BIOS is supplied by Digital Research, explicit instructions are provided for field reconfiguration of the BIOS to match nearly any hardware environment, see Section 6.

The BIOS and BDOS are logically combined into a single module with a common entry point and referred to as the FDOS. The CCP is a distinct program that uses the FDOS to provide a human-oriented interface with the information that is cataloged on the back-up and CCP, where various nonresident operating system commands and user programs are executed. The lower portion of memory is reserved for system information and is detailed in later sections. Memory organization of the CP/M system is shown in Figure 5-1.

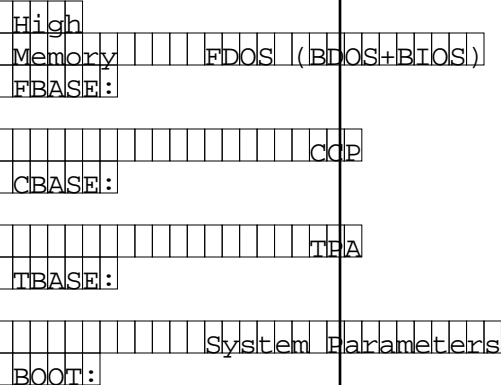


Figure 5-1. CP/M Memory Organization



The exact memory addresses corresponding to `BOOT`, `TBASE`, `CBASE`, and `FBASE` vary from version to version and are described fully in Section 6. All standard CP/M versions assume `BOOT=0000H`, which is the base of random access memory. The machine code found at location `BOOT` performs a system warm start, which loads and initializes the programs and variables necessary to return control to the CCP. Thus, transient programs need only jump to location `BOOT` to return control to CP/M at the command level. Further, the standard versions assume `TBASE=BOOT+0100H`, which is normally location `0100H`. The principal entry point to the FDOS is at location `BOOT+0005H` (normally `0005H`) where a jump to `FBASE` is found. The address field at `BOOT+0006H` (normally `0006H`) contains the value of `FBASE` and can be used to determine the size of available memory, assuming that the CCP is being overlaid by a transient program.

Transient programs are loaded into the TPA and executed as follows. The operator communicates with the CCP by typing command lines following each prompt. Each command line takes one of the following forms:

```
command
command file1
command file1 file2
```

where `command` is either a built-in function, such as `DIR` or `TYPE`, or the name of a transient command or program. If the command is a built-in function of CP/M, it is executed immediately. Otherwise,

```
command.COM
```

If the file is found, it is assumed to be a memory image of a program that executes in the TPA and thus implicitly originates at `TBASE` in memory. The CCP loads the `COM` file from the disk into memory starting at `TBASE` and can extend up to `CBASE`.

If the command is followed by one or two file specifications, the CCP prepares one or two File Control Block (FCB) names in the system parameter area. These optional FCBs are in the form necessary to access files through the FDOS and are described in Section 5.2.

The transient program receives control from the CCP and begins execution, using the I/O facilities of the FDOS. The transient program is called from the CCP. Thus, it can simply return to the CCP upon completion of its processing, or can jump to `BOOT` to pass control back to CP/M. In the first case, the transient program must not use memory above `CBASE`, while in the latter case, memory up through `FBASE-1` can be used.

The transient program can use the CP/M I/O facilities to communicate with the operator's console and peripheral devices, including the disk subsystem. The I/O system is accessed by passing a function number and an information address to CP/M through the FDOS entry point at `BOOT+0005H`. In the case of a disk read, for



## 5.2 Operating System Call Conventions

```
o read a console character
o write a console character
o read a sequential character
o write a sequential character
o get or set I/O status
o print console buffer
o interrogate console ready
```

- o disk system reset
- o drive selection
- o file creation
- o file close
- o directory search
- o file delete
- o file rename
- o random or sequential read
- o random or sequential write
- o interrogate available disks
- o interrogate selected disk
- o set DMA address
- o set/reset file indicators.

5	-	3
---	---	---



CP/M Operating System Manual 5.2 Call Conventions

0	System Reset	19	Delete File
1	Console Input	20	Read Sequential
2	Console Output	21	Write Sequential
3	Reader Input	22	Make File
4	Punch Output	23	Rename File
5	List Output	24	Return Login Vector
6	Direct Console I/O	25	Return Current Disk
7	Get I/O Byte	26	Set DMA Address
8	Set I/O Byte	27	Get Addr(Alloc)
9	Print String	28	Write Protect Disk
10	Read Console Buffer	29	Get R/O Vector
11	Get Console Status	30	Set File Attributes
12	Return Version Number	31	Get Addr(Disk Parms)
13	Reset Disk System	32	Set/Get User Code
14	Select Disk	33	Read Random
15	Open File	34	Write Random
16	Close File	35	Compute File Size
17	Search for First	36	Set Random Record
18	Search for Next	37	Reset Drive
		40	Write Random with Zero Fill

Functions 28 and 32 should be avoided in application programs to maintain upward compatibility with CP/M.

Upon entry to a transient program, the CCP leaves the stack pointer set to an eight-level stack area with the CCP return address

Although this stack is usually not used by a transient program (most transients return to the CCP through a jump to location 0000H) it is large enough to make CP/M system calls because the FDOS switches to a local stack at system entry. For example, the assembly-language program segment below reads characters continuously until an asterisk is encountered, at which time control returns to the CCP, assuming a standard CP/M system with BOOT = 0000H.

```
BDOS      EQU      0005H      ;STANDARD CP/M ENTRY
CONIN     EQU      1          ;CONSOLE INPUT FUNCTION
;
;
;
NEXTC:    ORG      0100H      ;BASE OF TPA
          MVI      C,CONIN    ;READ NEXT CHARACTER
          CALL     BDOS        ;RETURN CHARACTER IN <A>
          CPI      '*'        ;END OF PROCESSING?
          JNZ      NEXTC      ;LOOP IF NOT
          RET              ;RETURN TO CCP
          END
```

CP/M implements a named file structure on each disk, providing a logical organization that allows any particular file to contain any number of records from completely empty to the full capacity of the drive. Each drive is logically distinct with a disk directory and file data area. The disk filenames are in three parts: the drive select code, the filename (consisting of one to eight nonblank



characters), and the filetype (consisting of zero to three nonblank characters). The filetype names the generic category of a particular file, while the filename distinguishes individual files in each category. The filetypes listed in Table 5-1 name a few generic categories that have been established, although they are somewhat arbitrary.

Table 5-1. CP/M Filetypes

Filetype	Meaning
ASM	Assembler Source
PRN	Printer Listing
HEX	Hex Machine Code
BAS	Basic Source File
INT	Intermediate Code
COM	Command File
PLI	PL/I Source File
REL	Relocatable Module
TEX	TEX Formatter Source
BAK	ED Source Backup
SYM	SID Symbol File
\$\$\$	Temporary File

Source files are treated as a sequence of ASCII characters, and line-feed sequence (0DH followed by 0AH). Thus, one 128-byte CP/M record can contain several lines of source text. The end of an ASCII file is denoted by a CTRL-Z character (1AH) or a real end-of-file returned by the CP/M read operation. CTRL-Z characters embedded within machine code files (for example, COM files) are ignored and the end-of-file condition returned by CP/M is used to terminate read operations.

Files in CP/M can be thought of as a sequence of up to 65536 records of 128 bytes each, numbered from 0 through 65535, thus allowing a maximum of 8 megabytes per file. Note, however, that although the records may be considered logically contiguous, they may not be physically contiguous in the disk data area. Internally, all files are divided into 16K byte segments called logical extents, so that counters are easily maintained as 8-bit values. The division into extents is discussed in the paragraphs that follow; however, they are not particularly significant for the programmer, because each extent is automatically accessed in both sequential and random access modes.

In the file operations starting with Function 15, DE usually addresses a FCB. Transient programs often use the default FCB area reserved by CP/M at location BOOT+005CH (normally 005CH) for simple file operations. The basic unit of file information is a 128-byte record used for all file operations. Thus, a default location for disk I/O is provided by CP/M at location BOOT+0080H (normally 0080H) which is the initial default DMA address. See Function 26.



All directory operations take place in a reserved area that does not affect write buffers as was the case in release 1, with the exception of Search First and Search Next, where compatibility is required.

The FCB data area consists of a sequence of 33 bytes for sequential access and a series of 36 bytes in the case when the file is accessed randomly. The default FCB, normally located at 005CH, can be used for random access files, because the three bytes starting at BOOT+007DH are available for this purpose. Figure 5-2 shows the FCB format with the following fields.

dr	f1	f2	/	/	f8	t1	t2	t3	ex	s1	s2	rc	d0	/	/	dn	cr	r0	r1	r2
00	01	02	...	08	09	10	11	12	13	14	15	16	...	31	32	33	34	35		

Figure 5-2. File Control Block Format

The following table lists and describes each of the fields in the File Control Block figure.

Table 5-2. File Control Block Fields

Field	Definition
dr	drive code ((0-16) 0 = use default drive for file 1 = auto disk select drive A, 2 = auto disk select drive B, . . . 16= auto disk select drive P.
f1....f8	contain the filename in ASCII upper-case, with high bit = 0
t1, t2, t3	contain the filetype in ASCII upper-case, with high bit = 0 t1', t2', and t3' denote the bit of these positions, t1' = 1 =>Read-Only file, t2' = 1 =>SYS file, no DIR list
ex	contains the current extent number, normally set to 00 by the user, but in range 0-31 during file I/O



Table 5-2. (continued)

Field	Definition
s1	reserved for internal system use
s2	reserved for internal system use, set to zero on call to OPEN, MAKE, SEARCH
rc	record count for extent ex; takes on values from 0-127
d0....dn	filled in by CP/M; reserved for system use
cr	current record to read or write in a sequential file operation; normally set to zero by user
r0, r1, r2	optional random record number in the range 0-65535, with overflow to r2, r0, r1 constitute a 16-bit value with low byte r0, and high byte r1

FCB, which provides the name and allocation information for all subsequent file operations. When accessing files, it is the programmer's responsibility to fill the lower 16 bytes of the FCB and initialize the cr field. Normally, bytes 1 through 11 are set to the ASCII character values for the filename and filetype, while all other fields are zero.

FCBs are stored in a directory area of the disk, and are brought into central memory before the programmer proceeds with file operations (see the OPEN and MAKE functions). The memory copy of the FCB is updated as file operations take place and later recorded permanently on disk at the termination of the file operation, (see the CLOSE command).

The CCP constructs the first 16 bytes of two optional FCBs for a transient by scanning the remainder of the line following the transient name, denoted by file1 and file2 in the prototype command line described above, with unspecified fields set to ASCII blanks. The first FCB is constructed at location BOOT+005CH and can be used as is for subsequent file operations. The second FCB occupies the d0....dn portion of the first FCB and must be moved to another area of memory before use. If, for example, the following command line is typed:

```
PROGNAME B:X.ZOT Y.ZAP
```



## 5.2 Call Conventions

As an added convenience, the default buffer area at location `BOOT+0080H` is initialized to the command line tail typed by the operator following the program name. The first position contains the number of characters, with the characters themselves following the character count. Given the above command line, the area beginning at `BOOT+0080H` is initialized as follows:

	+00	+01	+02	+03	+04	+05	+06	+07	+08	+09	+A	+B	+C	+D	+E
--	-----	-----	-----	-----	-----	-----	-----	-----	-----	-----	----	----	----	----	----

Individual functions are described in detail in the pages that follow.



## 5.2 Call Conventions

Register 0:	00H
-------------	-----

5	-	9
---	---	---



## Call Conventions

## FUNCTION 2: CONSOLE OUTPUT

### Entry Parameters

Register C: 02H

Register E: ASCII Character

The ASCII character from register E is sent to the console device. As in Function 1, tabs are expanded and checks are made for start/stop scroll and printer echo.

### FUNCTION 3: READER INPUT

### Entry Parameters:

Register C: 03H

Returned Value:

The Reader Input function reads the next character from the logical reader into register A. See the IOBYTE definition in Chapter 6. Control does not return until the character has been read.	
--	--



## Call Conventions

## FUNCTION 4: PUNCH OUTPUT

Entry Parameters:

Register C:	04H
-------------	-----

```

register E:  ASCII Character

```

The Punch Output function sends the character from register E	
to the logical punch device.	

## FUNCTION 5: LIST OUTPUT

### Entry Parameters:

Register C:	05H
-------------	-----

Register E: ASCII Character

		The List Output function sends the ASCII character in register E to the logical listing device.	
--	--	---	--



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FUNCTION 6: DIRECT CONSOLE I/O

Entry Parameters:

Register C: 06H

Register E: 0FFH (input) or  
char (output)

Returned Value:

Register A: char or status

Direct Console I/O is supported under CP/M for those specialized applications where basic console input and output are required. Use of this function should, in general, be avoided since it bypasses all of the CP/M normal control character functions (for example, CTRL-S and CTRL-P). Programs that perform direct I/O through the BIOS under previous releases of CP/M, however, should be changed to use direct I/O under BDOS so that they can be fully supported under future releases of MP/M and CP/M.

Upon entry to Function 6, register E either contains hexadecimal FF, denoting a console input request, or an ASCII character. If E contains FF, register A contains the next console input character if no character is ready, otherwise A contains the next console input character.

If the input value in E is not FF, Function 6 assumes that E contains a valid ASCII character that is sent to the console.

Function 6 must not be used in conjunction with other console I/O functions.

FUNCTION 7: GET I/O BYTE

Entry Parameters:

Register C: 07H

Returned Value:

Register A: I/O Byte Value

The Get I/O Byte function returns the current value of IOBYTE in register A. See Chapter 6 for IOBYTE definition.



## 5.2 Call Conventions

5	-	1	3
---	---	---	---



FUNCTION 10: READ CONSOLE BUFFER

Entry Parameters:

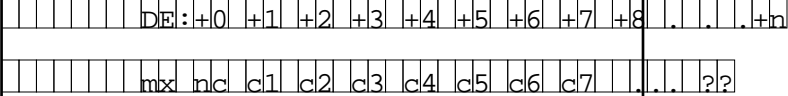
Register C: 0AH

Registers DE: Buffer Address

Returned Value:

Console Characters in Buffer

The Read Buffer functions reads a line of edited console input into a buffer addressed by registers DE. Console input is terminated when either input buffer overflows or a carriage return or line-feed is typed. The Read Buffer takes the form:



where mx is the maximum number of characters that the buffer will hold, 1 to 255, and nc is the number of characters read (set by FDOS upon return) followed by the characters read from the console. If denoted by ?? in the above figure. A number of control functions, summarized in Table 5-3, are recognized during line editing.

Table 5-3. Edit Control Characters

Character	Edit Control Function
rub/del	removes and echoes the last character
CTRL-C	reboots when at the beginning of line
CTRL-E	causes physical end of line
CTRL-H	backspaces one character position
CTRL-J	(line feed) terminates input line
CTRL-M	(return) terminates input line
CTRL-R	retypes the current line after new line
CTRL-U	removes current line
CTRL-X	same as CTRL-U







## 5.2 Call Conventions

5	-	1	6
---	---	---	---



## 5.2 Call Conventions

Register C:	0DH
-------------	-----

Register C: 0EH

5-17







## Call Conventions

FUNCTION 16: CLOSE FILE
-------------------------

### Entry Parameters:

Register C: 10H

Registers DE: FCB Address

Returned Value:
-----------------

Register A: Directory Code

The Close File function performs the inverse of the Open File function. Given that the FCB addressed by DE has been previously activated through an open or make function, the close function permanently records the new FCB in the reference disk directory see functions 15 and 22. The FCB matching process for the close is identical to the open function. The directory code returned for a successful close operation is 0, 1, 2, or 3, while a 0FFH (255 decimal) is returned if the filename cannot be found in the directory. A file need not be closed if only read operations have taken place. If write operations have occurred, the close operation is necessary to record the new directory information permanently.







## 5.2 Call Conventions

## FUNCTION 18: SEARCH FOR NEXT

### Entry Parameters:

Register C:	12H
-------------	-----

Returned Value:

Register A: Directory Code

The Search Next function is similar to the Search First function, except that the directory scan continues from the last matched entry. Similar to Function 17, Function 18 returns the decimal value 255 in A when no more directory items match.

FUNCTION	19:	DELETE	FILE
----------	-----	--------	------

Entry Parameters:
-------------------

Registers DE:	FCB Address
---------------	-------------

Returned Value:
-----------------

Register	A:	Directory Code
----------	----	----------------

The Delete File function removes files that match the FCB addressed by DE. The filename and type may contain ambiguous references (that is, question marks in various positions), but the drive select code cannot be ambiguous, as in the Search and Search Next functions.

Function 19 returns a decimal 255 if the referenced file or files cannot be found; otherwise, a value in the range 0 to 3 returned.
---



## 5.2 Call Conventions

5-22







Registers HL: Log-in Vector

Register	A:	Current Disk
----------	----	--------------

5-24



## Call Conventions

FUNCTION	26:	SET	DMA	ADDRESS
----------	-----	-----	-----	---------

### Entry Parameters:

```
Register C: 1AH
```

Registers DE:	DMA Address
---------------	-------------

DMA is an acronym for Direct Memory Address, which is often used in connection with disk controllers that directly access the memory of the mainframe computer to transfer data to and from the disk subsystem. Although many computer systems use non-DMA access (that is, the data is transferred through programmed I/O operations), the DMA address has, in CP/M, come to mean the address at which the 128-byte data record resides before a disk write and after a disk read. Upon cold start, warm start, or disk system reset, the DMA address is automatically set to BOOT+0080H. The Set DMA function can be used to change this default value to address another area of memory where the data records reside. Thus, the DMA address becomes the value specified by DE until it is changed by a subsequent Set DMA function, cold start, warm start, or disk system reset.

FUNCTION	27:	GET ADDR (ALLOC)
----------	-----	------------------

### Entry Parameters:

Register C:	1BH
-------------	-----

Returned Value:

Registers HL: ALLOC Address

An allocation vector is maintained in main memory for each on-line disk drive. Various system programs use the information provided by the allocation vector to determine the amount of remaining storage (see the STAT program). Function 27 returns the base address of the allocation vector for the currently selected disk drive. However, the allocation information might be invalid if the selected disk has been marked Read-Only. Although this function is not normally used by application programs, additional details of the allocation vector are found in Chapter 6.



## 5-26



## Call Conventions

5	-	2	7
---	---	---	---



## Call Conventions

An application program can change or interrogate the currently active user number by calling Function 32. If register E = 0FFH, the value of the current user number is returned in register A, where the value is in the range of 0 to 15. If register E is not 0FFH, the current user number is changed to the value of E, modulo 16.



## FUNCTION 33: READ RANDOM

## Entry Parameters:

Register C: 21H

## Returned Value:

Register A: Return Code

The Read Random function is similar to the sequential file read operation of previous releases, except that the read operation takes place at a particular record number, selected by the 24-bit value constructed from the 3-byte field following the FCB (byte positions r0 at 33, r1 at 34, and r2 at 35). The user should note that the sequence of 24 bits is stored with least significant byte first (r0), middle byte next (r1), and high byte last (r2). CP/M does not reference byte r2, except in computing the size of a file (Function 35). Byte r2 must be zero, however, since a nonzero value indicates overflow past the end of file.

Thus, the r0, r1 byte pair is treated as a double-byte, or word value, that contains the record to read. This value ranges from 0 megabyte file. To process a file using random access, the base extent (extent 0) must first be opened. Although the base extent might or might not contain any allocated data, this ensures that the file is properly recorded in the directory and is visible in DIR requests. The selected record number is then stored in the random record field (r0, r1), and the BDOS is called to read the record.

Upon return from the call, register A either contains an error code, as listed below, or the value 00, indicating the operation was successful. In the latter case, the current DMA address contains the randomly accessed record. Note that contrary to the sequential read operation, the record number is not advanced. Thus, subsequent random read operations continue to read the same record.

Upon each random read operation, the logical extent and current record values are automatically set. Thus, the file can be sequentially read or written, starting from the current randomly accessed position. However, note that, in this case, the last randomly read record will be reread as one switches from random mode to sequential read and the last record will be rewritten as one switches to a sequential write operation. The user can simply advance the random record position following each random read or write to obtain the effect of sequential I/O operation.



## Call Conventions

Error codes 01 and 04 occur when a random read operation accesses a data block that has not been previously written or an extent that has not been created, which are equivalent conditions. Error code 03 does not normally occur under proper system operation. If it does, it can be cleared by simply rereading or reopening extent zero as long as the disk is not physically write protected. Error code 06 occurs whenever byte r2 is nonzero under the current 2.0 release. Normally, nonzero return codes can be treated as missing data, with zero return codes indicating operation complete.











FUNCTION 36: SET RANDOM RECORD

Entry Parameters:

Register C: 24H

Registers DE: FCB Address

Returned Value:

Random Record Field Set

The Set Random Record function causes the BDOS automatically to produce the random record position from a file that has been read or written sequentially to a particular point. The function can be useful in two ways.

First, it is often necessary initially to read and scan a sequential file to extract the positions of various key fields. As each key is encountered, Function 36 is called to compute the random record position for the data corresponding to this key. If the data unit size is 128 bytes, the resulting record position is placed into a table with the key for later retrieval. After scanning the entire file and tabulating the keys and their record numbers, the user can

read, using the corresponding random record number that was saved earlier. The scheme is easily generalized for variable record lengths, because the program need only store the buffer-relative byte position along with the key and record number to find the exact starting position of the keyed data at a later time.

A second use of Function 36 occurs when switching from a sequential read or write over to random read or write. A file is sequentially accessed to a particular point in the file, Function 36 is called, which sets the record number, and subsequent random read and write operations continue from the selected point in the file.



## Call Conventions

Register	A:	00H
----------	----	-----

To maintain compatibility with MP/M, CP/M returns a zero value.

Register A: Return Code
-------------------------

5	-	3	4
---	---	---	---



[illegible]



5-36



			i	console messages
0187	6e6f20f		nofile:	db 'no source file\$'
0196	6e6f209		nodir:	db 'no directory space\$'
01a9	6f7574f		space:	db 'out of dat space\$'
01bb	7772695		wrprot:	db 'write protected?\$'
01cc	636f700		normal:	db 'copy complete\$'
			i	
			i	data areas
01da			dfcb:	ds 33 ;destination fcb
01fa			dfcbcr	equ dfcb+32 ;current record
			i	
01fb				ds 32 ;16 level stack
			stack:	
021b				end

5	-	3	7
---	---	---	---



## 5.4 A Sample File Dump Utility

x	.	i	n		5
---	---	---	---	--	---

 $i$ [illegible] $i$ 

```

; non graphic characters

```

```
000d = cr equ 0dh ;carriage return
```

```
000a = If equ 0ah ;line feed
```

*i*

```
; file control block definitions
```

005c	=	fcldn	egu	fcld+0	;disk name
------	---	-------	-----	--------	------------

```
005d = fcbfn equ fcb+1 ;file name
```

0065 =	fcbf	equ	fcbf+9	disk file type (3 characters)
--------	------	-----	--------	----------------------------------

0068	=	fcbrl	equ	fcb+12	;file's current reel ;number
------	---	-------	-----	--------	---------------------------------

006b =	fcbrc	equ	fcbr+15	;file's record count (0 to	
				;128)128)	

```
007c = fcbcr' equ fcb+32 ;current (next) record
```







[illegible]



```
ret
```

7800012305	CP1	8011
------------	-----	------

```
read another buffer
```

SIC
ret

```

;index to de

```

```
sta    ibp        ;back to memory
```

```
save the current file address
```

dad	d
-----	---

dad	d
chre	ute
che	meters
e	dd
na	g
i	c
i	n
b	a

byte	is	in	the	accumulator
------	----	----	-----	-------------

```
ret
```

```
i set up file
```















[illegible]



[illegible]







[illegible]



Major improvements could be made to this particular program to enhance its operation. In fact, with some work, this program could evolve into a simple data base management system. One could, for example, assume a standard record size of 128 bytes, consisting to arbitrary fields within the record. A program, called GETKEY, could be developed that first reads a sequential file and extracts a specific field defined by the operator. For example, the command

```
GETKEY NAMES.DAT LASTNAME 10 20
```

would cause GETKEY to read the data base file NAMES.DAT and extract the LAST-NAME field from each record, starting in position 10 and ending at character 20. GETKEY builds a table in memory consisting of each particular LASTNAME field, along with its 16-bit record number location within the file. The GETKEY program then sorts this list and writes a new file, called LASTNAME.KEY, which is an alphabetical list of LASTNAME fields with their corresponding record numbers. This list is called an inverted index in information retrieval parlance.

If the programmer were to rename the program shown above as QUERY and modify it so that it reads a sorted key file into memory, the command line might appear as

```
QUERY NAMES.DAT LASTNAME.KEY
```

string that is a particular key to find in the NAMES.DAT data base. Because the LASTNAME.KEY list is sorted, one can find a particular entry rapidly by performing a binary search, similar to looking up a name in the telephone book. Starting at both ends of the list, one examines the entry halfway in between and, if not matched, splits either the upper half or the lower half for the next search. You will quickly reach the item you are looking for and find the corresponding record number. You should fetch and display this record at the console, just as was done in the program shown above.

With some more work, you can allow a fixed grouping size that differs from the 128-byte record shown above. This is accomplished by keeping track of the record number and the byte offset within the record. Knowing the group size, you randomly access the record containing the proper group, offset to the beginning of the group within the record read sequentially until the group size has been exhausted.

Finally, you can improve QUERY considerably by allowing boolean expressions, which compute the set of records that satisfy several relationships, such as a LASTNAME between HARDY and LAUREL and an AGE lower than 45. Display all the records that fit this description. Finally, if your lists are getting too big to fit into memory, randomly access key files from the disk as well.











## Section 6

### CP/M 2 Alteration

#### 6.1 Introduction

The standard CP/M system assumes operation on an Intel Model 800 microcomputer development system, but is designed so you can alter a specific set of subroutines that define the hardware operating environment.

Although standard CP/M 2 is configured for single-density floppy disks, field-alteration features allow adaptation to a wide variety of disk subsystems from single-drive minidisks to high-capacity, hard disk systems. To simplify the following adaptation process, it is assumed that CP/M 2 is first configured for single-density floppy disks where minimal editing and debugging tools are available. If an earlier version of CP/M is available, the customizing process is eased considerably. In this latter case, you might want to review the system generation process and skip to later sections that discuss system alteration for nonstandard disk systems.

To achieve device independence, CP/M is separated into three distinct modules:

- o BIOS is the Basic I/O System, which is environment dependent.
- o BDOS is the Basic Disk Operating System, which is not dependent upon the hardware configuration.
- o CCP is the Console Command Processor, which uses the BDOS.

Of these modules, only the BIOS is dependent upon the particular hardware. You can patch the distribution version of CP/M to provide a new BIOS that provides a customized interface between the remaining CP/M modules and the hardware system. This document provides a step-by-step procedure for patching a new BIOS into CP/M.

All disk-dependent portions of CP/M 2 are placed into a BIOS, a resident disk parameter block, which is either hand coded or produced automatically using the disk definition macro library provided with CP/M 2. The end user need only specify the maximum number of active disks, the starting and ending sector numbers, the data allocation size, the maximum extent of the logical disk, directory size information, and reserved track values. The macros use this information to generate the appropriate tables and table references for use during CP/M 2 operation. Deblocking information is provided, which aids in assembly or disassembly of sector sizes that are multiples of the fundamental 128-byte data unit, and the system alteration manual includes general purpose subroutines that use the deblocking information to take advantage of larger sector sizes. Use of these subroutines, together with the table-drive data access algorithms, makes CP/M 2 a universal data management system.



File expansion is achieved by providing up to 512 logical file extents, where each logical extent contains 16K bytes of data. CP/M 2 is structured, however, so that as much as 128K bytes of data are addressed by a single physical extent, corresponding to a single directory entry, maintaining compatibility with previous versions while taking advantage of directory space.

If CP/M is being tailored to a computer system for the first time, the new BIOS requires some simple software development and testing. The standard BIOS is listed in Appendix A and can be used as a model for the customized package. A skeletal version of the BIOS given in Appendix B can serve as the basis for a modified BIOS.

In addition to the BIOS, you must write a simple memory loader, called GETSYS, which brings the operating system into memory. To patch the new BIOS into CP/M, you must write the reverse of GETSYS, called PUTSYS, which places an altered version of CP/M back onto the disk. PUTSYS can be derived from GETSYS by changing the disk read commands into disk write commands. Sample skeletal GETSYS and PUTSYS programs are described in Section 6.4 and listed in Appendix C.

To make the CP/M system load automatically, you must also supply a cold start loader, similar to the one provided with CP/M, listed in Appendixes A and D. A skeletal form of a cold start loader is given in Appendix E, which serves as a model for the loader.

## 6.2 First-level System Regeneration

The procedure to patch the CP/M system is given below. Address references in each step are shown with H denoting the hexadecimal radix, and are given for a 20K CP/M system. For larger CP/M systems, a bias is added to each address that is shown with a +b following it, where b is equal to the memory size-20K. Values for b in various standard memory sizes are listed in Table 6-1.

Table 6-1. Standard Memory Size Values

Memory Size	Value
24K:            b = 24K - 20K = 4K = 1000H	
32K:            b = 32K - 20K = 12K = 3000H	
40K:            b = 40K - 20K = 20K = 5000H	
48K:            b = 48K - 20K = 28K = 7000H	
56K:            b = 56K - 20K = 36K = 9000H	
62K:            b = 62K - 20K = 42K = A800H	
64K:            b = 64K - 20K = 44K = B000H	



Note that the standard distribution version of CP/M is set for operation within a 20K CP/M system. Therefore, you must first bring up the 20K CP/M system, then configure it for actual memory size (see Section 6.3).

Follow these steps to patch your CP/M system:

1) Read Section 6.4 and write a GETSYS program that reads the first two tracks of a disk into memory. The program from the disk must be loaded starting at location 3380H. GETSYS is coded to start at location 100H (base of the TPA) as shown in Appendix C.

2) Test the GETSYS program by reading a blank disk into memory, and check to see that the data has been read properly and that the disk has not been altered in any way by the GETSYS program.

3) Run the GETSYS program using an initialized CP/M disk to see if GETSYS loads CP/M starting at 3380H (the operating system actually starts 128 bytes later at 3400H).

4) Read Section 6.4 and write the PUTSYS program. This writes memory starting at 3380H back onto the first two tracks of the disk. The PUTSYS program should be located at 200H, as shown in Appendix C.

5) Test the PUTSYS program using a blank, uninitialized disk by writing a portion of memory to the first two tracks; clear memory and read it back using GETSYS. Test PUTSYS completely, because this program will be used to alter CP/M on disk.

6) Study Sections 6.5, 6.6, and 6.7 along with the distribution version of the BIOS given in Appendix A and write a simple version that performs a similar function for the customized environment. Use the program given in Appendix B as a model. Call this new BIOS by name CBIOS (customized BIOS). Implement only the primitive disk operations on a single drive and simple console input/output functions in this phase.

7) Test CBIOS completely to ensure that it properly performs console character I/O and disk reads and writes. Be careful to ensure that no disk write operations occur during read operations and check that the proper track and sectors are addressed on all reads and writes. Failure to make these checks might cause destruction of the initialized CP/M system after it is patched.

8) Referring to Table 6-3 in Section 6.5, note that the BIOS is placed between locations 4A00H and 4FFFH. Read the CP/M system using GETSYS and replace the BIOS segment by the CBIOS developed in step 6 and tested in step 7. This replacement is done in memory.







16) At this point, there is probably a good version of the customized CP/M system on the test disk. Use GETSYS to load CP/M from the test disk. Remove the test disk, place the distribution disk, or a legal copy, into the drive, and use PUTSYS to replace the distribution version with the customized version. Do not make this replacement if you are unsure of the patch because this step destroys the system that was obtained from Digital Research.

17) Load the modified CP/M system and test it by typing

DIR

CP/M responds with a list of files that are provided on the initialized disk. The file DDT.COM is the memory image for the debugger. Note that from now on, you must always reboot the CP/M system (CTRL-C is sufficient) when the disk is removed and replaced by another disk, unless the new disk is to be Read-Only.

18) Load and test the debugger by typing

DDT

See Chapter 4 for operating procedures.

19) Before making further CBIOS modifications, practice using

Recode and test the GETSYS, PUTSYS, and CBIOS programs using ED, ASM, and DDT. Code and test a COPY program that does a sector-to-sector copy from one disk to another to obtain back-up copies of the original disk. Read the CP/M Licensing Agreement specifying legal responsibilities when copying the CP/M system. Place the following copyright notice:

Copyright (c), 1983  
Digital Research

on each copy that is made with the COPY program.

20) Modify the CBIOS to include the extra functions for punches, readers, and sign-on messages, and add the facilities for additional disk drives, if desired. These changes can be made with the GETSYS and PUTSYS programs or by referring to the regeneration process in Section 6.3.

You should now have a good copy of the customized CP/M system. Although the CBIOS portion of CP/M belongs to the user, the modified version cannot be legally copied.

It should be noted that the system remains file-compatible with all other CP/M systems (assuming media compatibility) which allows transfer of nonproprietary software between CP/M users.



CP/M Operating System Manual 6.3 Second-level System Generation

### 6.3 Second-level System Generation

Once the system is running, the next step is to configure CP/M for the desired memory size. Usually, a memory image is first produced with the MOVCPM program (system relocater) and then placed into a named disk file. The disk file can then be loaded, examined, patched, and replaced using the debugger and the system generation program (refer to Chapter 1).

The CBIOS and BOOT are modified using ED and assembled using ASM, producing files called CBIOS.HEX and BOOT.HEX, which contain the code for CBIOS and BOOT in Intel hex format.

To get the memory image of CP/M into the TPA configured for the desired memory size, type the command:

[illegible]

where xx is the memory size in decimal K bytes, for example, 32 for 32K. The response is as follows:

```

CONSTRUCTING xxxK CP/M VERS 2.0

```

[illegible]

		"SAVE	34	CPMxx.COM"
--	--	-------	----	------------

memory size. The memory image is at location 0900H through 227FH, that is, the BOOT is at 0900H, the CCP is at 980H, the BDOS starts at 1180H, and the BIOS is at 1F80H. Note that the memory image has the standard Model 800 BIOS and BOOT on it. It is now necessary to save the memory image in a file so that you can patch the CBIOS and CBOOT into it:

[illegible]

The memory image created by the MOVCPM program is offset by a negative bias so that it loads into the free area of the TPA, and thus does not interfere with the operation of CP/M in higher memory. This memory image can be subsequently loaded under DDT and examined or changed in preparation for a new generation of the system. DDT is loaded with the memory image by typing:

```
DDT CPMxx.COM      Loads DDT, then reads the CP/M image.
```

DDT should respond with the following:

					NEXT	PC
					2300	0100

-	The DDT prompt
---	----------------

You can then give the display and disassembly commands to examine portions of the memory image between 900H and 227FH. Note, however, that to find any particular address within the memory image, you



must apply the negative bias to the CP/M address to find the actual address. Track 00, sector 01, is loaded to location 900H (the user should find the cold start loader at 900H to 97FH); track 00, sector 02, is loaded into 980H (this is the base of the CCP); and so on through the entire CP/M system load. In a 20K system, for example, the CCP resides at the CP/M address 3400H, but is placed into memory at 980H by the SYSGEN program. Thus, the negative bias, denoted by  $n$ , satisfies

$$3400H + n = 980H, \text{ or } n = 980H - 3400H$$

Assuming two's complement arithmetic,  $n = D580H$ , which can be checked by

$$3400H + D580H = 10980H = 0980H \text{ (ignoring high-order overflow).}$$

Note that for larger systems,  $n$  satisfies

$$\begin{aligned} (3400H + b) + n &= 980H, \text{ or} \\ n &= 980H - (3400H + b), \text{ or} \\ n &= D580H - b \end{aligned}$$

The value of  $n$  for common CP/M systems is given below.

**Table 6-2. Common Values for CP/M Systems**

Memory Size	BIAS $b$	Negative Offset $n$
20K	0000H	$D580H - 0000H = D580H$
24K	1000H	$D580H - 1000H = C580H$
32K	3000H	$D580H - 3000H = A580H$
40K	5000H	$D580H - 5000H = 8580H$
48K	7000H	$D580H - 7000H = 6580H$
56K	9000H	$D580H - 9000H = 4580H$
62K	A800H	$D580H - A800H = 2D80H$
64K	B000H	$D580H - B000H = 2580H$

If you want to locate the address  $x$  within the memory image loaded under DDT in a 20K system, first type

$Hx, n$  Hexadecimal sum and difference

and DDT responds with the value of  $x+n$  (sum) and  $x-n$  (difference). The first number printed by DDT is the actual memory address in the image where the data or code is located. For example, the following DDT command:

H3400, D580

produces 980H as the sum, which is where the CCP is located in the memory image under DDT.



[illegible]



Place the scratch disk in drive A, then perform a cold start to

6-9



[illegible]



.....  
perform disk read at this point, branch to  
label START if an error occurs  
.....

POP H ;RECOVER HL  
POP B ;RECOVER B AND C REGISTERS  
RET ;BACK TO MAIN PROGRAM

END START

#### Listing 6-1. (continued)

This program is assembled and listed in Appendix B for reference purposes, with an assumed origin of 100H. The hexadecimal operation codes that are listed on the left might be useful if the program has to be entered through the panel switches.

The PUTSYS program can be constructed from GETSYS by changing only a few operations in the GETSYS program given above, as shown in Appendix C. The register pair HL becomes the dump address, next address to write, and operations on these registers do not change within the program. The READSEC subroutine is replaced by a WRITESEC subroutine, which performs the opposite function: data from address HL is written to the track given by register B and sector  
PUTSYS into a single program during the test and development phase, as shown in Appendix C.

### 6.5 Disk Organization

The sector allocation for the standard distribution version of CP/M is given here for reference purposes. The first sector contains an optional software boot section (see the table on the following page. Disk controllers are often set up to bring track 0, sector 1, into memory at a specific location, often location 0000H. The program in this sector, called BOOT, has the responsibility of bringing the remaining sectors into memory starting at location 3400H+b. If the controller does not have a built-in sector load, the program in track 0, sector 1 can be ignored. In this case, load the program from track 0, sector 2, to location 3400H+b.

As an example, the Intel Model 800 hardware cold start loader brings track 0, sector 1, into absolute address 3000H. Upon loading this sector, control transfers to location 3000H, where the bootstrap operation commences by loading the remainder of track 0 and all of track 1 into memory, starting at 3400H+b. Note that this bootstrap loader is of little use in a non-microcomputer development system environment, although it is useful to examine it because some of the boot actions will have to be duplicated in the user's cold start loader.











Each jump address corresponds to a particular subroutine that performs the specific function, as outlined below. There are three major divisions in the jump table: the system reinitialization, which results from calls on BOOT and WBOOT; simple character I/O, performed by calls on CONST, CONIN, CONOUT, LIST, PUNCH, READER, and LISTST; and disk I/O, performed by calls on HOME, SETDSK, SETTRK, SETSEC, SETDMA, READ, WRITE, and SECTRAN.

All simple character I/O operations are assumed to be performed in ASCII, upper- and lower-case, with high-order (parity bit) set to zero. An end-of-file condition for an input device is given by an ASCII CTRL-Z (1AH). Peripheral devices are seen by CP/M as logical devices and are assigned to physical devices within the BIOS.

To operate, the BDOS needs only the CONST, CONIN, and CONOUT subroutines. LIST, PUNCH, and READER can be used by PIP, but not the BDOS. Further, the LISTST entry is currently used only by DESPOOL, the print spooling utility. Thus, the initial version of CBIOS can have empty subroutines for the remaining ASCII devices.

The following list describes the characteristics of each device.

- o CONSOLE is the principal interactive console that communicates with the operator and it is accessed through CONST, CONIN, and CONOUT. Typically, the CONSOLE is a device such as a CRT or

- o LIST is the principal listing device. If it exists on the user's system, it is usually a hard-copy device, such as a printer or teletype.

- o PUNCH is the principal tape punching device. If it exists, it is normally a high-speed paper tape punch or teletype.

- o READER is the principal tape reading device, such as a simple optical reader or teletype.

A single peripheral can be assigned as the LIST, PUNCH, and READER device simultaneously. If no peripheral device is assigned as the LIST, PUNCH, or READER device, the CBIOS gives an appropriate error message so that the system does not hang if the device is accessed by PIP or some other user program. Alternately, the PUNCH and LIST routines can just simply return, and the READER routine can return with a 1AH (CTRL-Z) in register A to indicate immediate end-of-file.

For added flexibility, you can optionally implement the IOBYTE function, which allows reassignment of physical devices. The IOBYTE function creates a mapping of logical-to-physical devices that can be altered during CP/M processing, see the STAT command in Section 1.6.1.







The implementation of the IOBYTE is optional and effects only the organization of the CBIOS. No CP/M systems use the IOBYTE (although they tolerate the existence of the IOBYTE at location 0003H) except for PIP, which allows access to the physical devices, and STAT, which allows logical-physical assignments to be make or displayed. For more information see Section 1. In any case the IOBYTE implementation should be omitted until the basic CBIOS is fully implemented and tested; then you should add the IOBYTE to increase the facilities.

Disk I/O is always performed through a sequence of calls on the various disk access subroutines that set up the disk number to access, the track and sector on a particular disk, and the Direct Memory Access (DMA) address involved in the I/O operation. After all these parameters have been set up, a call is made to the READ or WRITE function to perform the actual I/O operation.

There is often a single call to SETDISK to select a disk drive, followed by a number of read or write operations to the selected disk before selecting another drive for subsequent operations. Similarly, there might be a single call to set the DMA address, followed by several calls that read or write from the selected DMA address before the DMA address is changed. The track and sector subroutines are always called before the READ or WRITE operations are performed.

The READ and WRITE routines should perform several retries (10 the error condition is returned to the BDOS, it reports the error to the user. The HOME subroutine might or might not actually perform the track 00 seek, depending upon controller characteristics; the important point is that track 00 has been selected for the next operation and is often treated in exactly the same manner as SETTRK with a parameter of 00.

The following table describes the exact responsibilities of each BIOS entry point subroutine.

**Table 6-5. BIOS Entry Points**

Entry Point	Function
BOOT	The BOOT entry point gets control from the cold start loader and is responsible for basic system initialization, including sending a sign-on message, which can be omitted in the first version. If the IOBYTE function is implemented, it must be set at this point. The various system parameters that are set by the WBOOT entry point must be initialized, and control is transferred to the CCP at 3400+b for further processing. Note that register C must be set to zero to select drive A.



Table 6-5. (continued)

Entry Point	Function
-------------	----------

WBOOT	The WBOOT entry point gets control when a warm start occurs. A warm start is performed whenever a user program branches to location 0000H, or when the CPU is reset from the front panel. The CP/M system must be loaded from the first two tracks of drive A up to, but not including, the BIOS, or CBIOS, if the user has completed the patch. System parameters must be initialized as follows:
-------	--

location 0,1,2	Set to JMP WBOOT for warm starts (0000H: JMP 4A03H+b)
----------------	---

location 3	Set initial value of IOBYTE, if implemented in the CBIOS
------------	--

location 4	High nibble = current user no; low nibble = current drive
------------	--

location 5,6,7	Set to JMP BDOS, which is the primary entry point to CP/M for transient programs. (0005H: JMP 3C06H+b)
----------------	---

Refer to Section 6.9 for complete details of page zero use. Upon completion of the initialization, the WBOOT program must branch to the CCP at 3400H+b to restart the system. Upon entry to the CCP, register C is set to the drive to select after system initialization. The WBOOT routine should read location 4 in memory, verify that is a legal drive, and pass it to the CCP in register C.

CONST	You should sample the status of the currently assigned console device and return 0FFH in register A if a character is ready to read and 00H in register A if no console characters are ready.
-------	---

CONIN	The next console character is read into register A, and the parity bit is set, high-order bit, to zero. If no console character is ready, wait until a character is typed before returning.
-------	---



Table 6-5. (continued)

Entry Point	Function
-------------	----------

CONOUT	The character is sent from register C to the console output device. The character is in ASCII, with high-order parity bit set to zero. You might want to include a time-out on a line-feed or carriage return, if the console device requires some time interval at the end of the line (such as a TI Silent 700 terminal). You can filter out control characters that cause the console device to react in a strange way (CTRL-Z causes the Lear-Seigler terminal to clear the screen, for example).
--------	---

LIST	The character is sent from register C to the currently assigned listing device. The character is in ASCII with zero parity bit.
------	---

PUNCH	The character is sent from register C to the currently assigned punch device. The character is in ASCII with zero parity.
-------	---

READER	The next character is read from the currently assigned reader device into register A with zero parity (high-order bit must be zero); an ASCII CTRL-Z (1AH).
--------	---

HOME	The disk head of the currently selected disk (initially disk A) is moved to the track 00 position. If the controller allows access to the track 0 flag from the drive, the head is stepped until the track 0 flag is detected. If the controller does not support this feature, the HOME call is translated into a call to SETTRK with a parameter of 0.
------	--

SELDISK	The disk drive given by register C is selected for further operations, where register C contains 0 for drive A, 1 for drive B, and so on up to 15 for drive P (the standard CP/M distribution version supports four drives). On each disk select, SELDISK must return in HL the base address of a 16-byte area, called the Disk Parameter Header, described in Section 6.10. For standard floppy disk drives, the contents of the header and associated tables do not change; thus, the program segment included in the sample CBIOS performs this operation automatically.
---------	---



Table 6-5. (continued)

Entry Point	Function
-------------	----------

If there is an attempt to select a nonexistent drive, SETDSK returns HL=0000H as an error indicator. Although SETDSK must return the header address on each call, it is advisable to postpone the physical disk select operation until an I/O function (seek, read, or write) is actually performed, because disk selects often occur without ultimately performing any disk I/O, and many controllers unload the head of the current disk before selecting the new drive. This causes an excessive amount of noise and disk wear. The least significant bit of register E is zero if this is the first occurrence of the drive select since the last cold or warm start.

SETTRK	Register BC contains the track number for subsequent disk accesses on the currently selected drive. The sector number in BC is the same as the number returned from the SECTRAN entry point. You can choose to seek the selected track at this time or delay the seek until the next read or write actually occurs.
--------	---

76 corresponding to valid track numbers for standard floppy disk drives and 0-65535 for nonstandard disk subsystems.

SETSEC	Register BC contains the sector number, 1 through 26, for subsequent disk accesses on the currently selected drive. The sector number in BC is the same as the number returned from the SECTRAN entry point. You can choose to send this information to the controller at this point or delay sector selection until a read or write operation occurs.
--------	--

SETDMA	Register BC contains the DMA (Disk Memory Access) address for subsequent read or write operations. For example, if B = 00H and C = 80H when SETDMA is called, all subsequent read operations read their data into 80H through 0FFH and all subsequent write operations get their data from 80H through 0FFH, until the next call to SETDMA occurs. The initial DMA address is assumed to be 80H. The controller need not actually support Direct Memory Access. If, for example, all data transfers are through I/O ports, the CBIOS that is constructed uses the 128-byte area starting at the selected DMA address for the memory buffer during the subsequent read or write operations.
--------	--



Table 6-5. (continued)

Entry Point	Function
-------------	----------

READ	Assuming the drive has been selected, the track has been set, and the DMA address has been specified, the READ subroutine attempts to read one sector based upon these parameters and returns the following error codes in register A:
------	--

0	no errors occurred
---	--------------------

1	nonrecoverable error condition occurred
---	---

Currently, CP/M responds only to a zero or nonzero value as the return code. That is, if the value in register A is 0, CP/M assumes that the disk operation was completed properly. If an error occurs the CBIOS should attempt at least 10 retries to see if the error is recoverable. When an error is reported the BDOS prints the message BDOS ERR ONx: BAD SECTOR. The operator then has the option of pressing a carriage return to ignore the error, or CTRL-C to abort.

address to the currently selected drive, track, and sector. For floppy disks, the data should be marked as nondeleted data to maintain compatibility with other CP/M systems. The error codes given in the READ command are returned in register A, with error recovery attempts as described above.

LISTST	You return the ready status of the list device used by the DESPOOL program to improve console response during its operation. The value 00 is returned in A if the list device is not ready to accept a character and 0FFH if a character can be sent to the printer. A 00 value should be returned if LIST status is not implemented.
--------	---



Table 6-5. (continued)

Entry Point	Function
-------------	----------

SECTTRAN	Logical-to-physical sector translation is performed to improve the overall response of CP/M. Standard CP/M systems are shipped with a skew factor of 6, where six physical sectors are skipped between each logical read operation. This skew factor allows enough time between sectors for most programs to load their buffers without missing the next sector. In particular computer systems that use fast processors, memory, and disk subsystems, the skew factor might be changed to improve overall response. However, the user should maintain a single-density IBM-compatible version of CP/M for information transfer into and out of the computer system, using a skew factor of 6.
----------	--

	In general, SECTTRAN receives a logical sector number relative to zero in BC and a translate table address in DE. The sector number is used as an index into the translate table, with the resulting physical sector number in HL. For standard systems, the table and indexing code is provided in the CBIOS and need not be
--	---

## 6.7 A Sample BIOS

The program shown in Appendix B can serve as a basis for your first BIOS. The simplest functions are assumed in this BIOS, so that you can enter it through a front panel, if absolutely necessary. You must alter and insert code into the subroutines for CONST, CONIN, CONOUT, READ, WRITE, and WAITIO subroutines. Storage is reserved for user-supplied code in these regions. The scratch area reserved in page zero (see Section 6.9) for the BIOS is used in this program, so that it could be implemented in ROM, if desired.

Once operational, this skeletal version can be enhanced to print the initial sign-on message and perform better error recovery. The subroutines for LIST, PUNCH, and READER can be filled out and the IOBYTE function can be implemented.

## 6.8 A Sample Cold Start Loader

The program shown in Appendix E can serve as a basis for a cold start loader. The disk read function must be supplied by the user, and the program must be loaded somehow starting at location 0000. Space is reserved for the patch code so that the total amount of storage required for the cold start loader is 128 bytes.



CP/M Operating System Manual 6.8 A Sample Cold Start Loader

Eventually, you might want to get this loader onto the first disk sector (track 0, sector 1) and cause the controller to load it into memory automatically upon system start up. Alternatively, the cold start loader can be placed into ROM, and above the CP/M system. In this case, it is necessary to originate the program at a higher address and key in a jump instruction at system start up that branches to the loader. Subsequent warm starts do not require this key-in operation, because the entry point WBOOT gets control, thus bringing the system in from disk automatically. The skeletal cold start loader has minimal error recovery, which might be enhanced in later versions.

## 6.9 Reserved Locations in Page Zero

Main memory page zero, between locations 00H and 0FFH, contains several segments of code and data that are used during CP/M processing. The code and data areas are given in the following table.

Table 6-6. Reserved Locations in Page Zero

Locations	Contents
000H-0002H	Contains a jump instruction to the warm start entry location 4A03H+b. This (0000H) or manual restart from the front panel.
0003H-0003H	Contains the Intel standard IOBYTE, is optionally included in the user's CBIOS (refer to Section 6.6).
0004H-0004H	Current default drive number (0=A, ..., 15=P).
0005H-0007H	Contains a jump instruction to the BDOS and serves two purposes: JMP 0005H provides the primary entry point to the BDOS, as described in Chapter 5, and LHLD 0006H brings the address field of the instruction to the HL register pair. This value is the lowest address in memory used by CP/M, assuming the CCP is being overlaid. The DDT program changes the address field to reflect the reduced memory size in debug mode.
0008H-0027H	Interrupt locations 1 through 5 not used.
0030H-0037H	Interrupt location 6 (not currently used) is reserved.



Table 6-6. (continued)

Locations	Contents
0038H-003AH	Restart 7; contains a jump instruction into the DDT or SID program when running in debug mode for programmed breakpoints, but is not otherwise used by CP/M.
003BH-003FH	Not currently used; reserved.
0040H-004FH	A 16-byte area reserved for scratch by CBIOS, but is not used for any purpose in the distribution version of CP/M.
0050H-005BH	Not currently used; reserved.
005CH-007CH	Default File Control Block produced for a transient program by the CCP.
007DH-007FH	Optional default random record position.
0080H-00FFH	Default 128-byte disk buffer, also filled with the command line when a transient is loaded under the CCP.

This information is set up for normal operation under the CP/M system, but can be overwritten by a transient program if the BDOS facilities are not required by the transient.

If, for example, a particular program performs only simple I/O and must begin execution at location 0, it can first be loaded into the TPA, using normal CP/M facilities, with a small memory move program that gets control when loaded. The memory move program must get control from location 0100H, which is the assumed beginning of all transient programs. The move program can then proceed to the entire memory image down to location 0 and pass control to the starting address of the memory load.

If the BIOS is overwritten or if location 0, containing the warm start entry point, is overwritten, the operator must bring the CP/M system back into memory with a cold start sequence.

## 6.10 Disk Parameter Tables

Tables are included in the BIOS that describe the particular characteristics of the disk subsystem used with CP/M. These tables can be either hand-coded, as shown in the sample CBIOS in Appendix B, or automatically generated using the DISKDEF macro library, as shown in Appendix F. The purpose here is to describe the elements of these tables.



CP/M Operating System Manual															
6.10 Disk Parameter Tables															
<p>In general, each disk drive has an associated (16-byte) disk parameter header that contains information about the disk drive and provides a scratch pad area for certain BDOS operations. The format of the disk parameter header for each drive is shown in Figure 6-2, where each element is a word (16-bit) value.</p>															
XLT	0000	0000	0000	DIRBUF	DPB	CSV	ALV								
16b	16b	16b	16b	16b	16b	16b	16b								

Figure 6-2. Disk Parameter Header Format

The meaning of each Disk Parameter Header (DPH) element is detailed in Table 6-7.

Table 6-7. Disk Parameter Headers

Disk Parameter Header				Meaning											
XLT				Address of the logical-to-physical particular drive, or the value 0000H if no sector translation takes place (that is, the physical and logical sector numbers are the same). Disk drives with identical sector skew factors share the same translate tables.											
0000				Scratch pad values for use within the BDOS, initial value is unimportant.											
DIRBUF				Address of a 128-byte scratch pad area for directory operations within BDOS. All DPHs address the same scratch pad area.											
DPB				Address of a disk parameter block for this drive. Drives with identical disk characteristics address the same disk parameter block.											
CSV				Address of a scratch pad area used for software check for changed disks. This address is different for each DPH.											
ALV				Address of a scratch pad area used by the BDOS to keep disk storage allocation information. This address is different for each DPH.											



CP/M Operating System Manual 6.10 Disk Parameter Tables

Given  $n$  disk drives, the DPHs are arranged in a table whose first row of 16 bytes corresponds to drive 0, with the last row corresponding to drive  $n-1$ . In the following figure the table DPBASE defines the base address of the DPH table.

DPBASE:

00	XLT	00	0000	0000	0000	DIRBUF	DBP	00	CSV	00	ATV	00
01	XLT	01	0000	0000	0000	DIRBUF	DBP	01	CSV	01	ATV	01
.	.	.	.	.	.	.	.	.	.	.	.	.
n-1	XLT	n-1	0000	0000	0000	DIRBUF	DBT	n-1	CSV	n-1	ATV	n-1

Figure 6-3. Disk Parameter Header Table

A responsibility of the SEIDSK subroutine is to return the base address of the DPH for the selected drive. The following sequence of operations returns the table address, with a 0000H returned if the selected drive does not exist.

```
NDISKS EQU 4 ;NUMBER OF DISK DRIVES
*****
SEIDSK: ;SELECT DISK GIVEN BY BC
        LSI H,0000H ;ERROR CODE
        MOV A,C ;DRIVE OK?
        CPI NDISKS ;CY IF SO
        RNC ;RET IF ERROR
        ;NO ERROR, CONTINUE
        MOV L,C ;LOW(DISK)
        MOV H,B ;HIGH(DISK)
        DAD H ;*2
        DAD H ;*4
        DAD H ;*8
        DAD H ;*16
        LXI D,DPBASE;FIRST DPH
        DAD D ;DPH(DISK)
        RET
```

The translation vectors, XLT 00 through XLT  $n-1$ , are located elsewhere in the BIOS, and simply correspond one-for-one with the logical sector numbers zero through the sector count 1. The Disk Parameter Block (DPB) for each drive is more complex. As shown in Figure 6-4, particular DPB, that is addressed by one or more DPHs, takes the general form:



CP/M Operating System Manual										6.10 Disk Parameter Tables									
	SPT	BSH	BLM	EXM	DSM	DRM	AL0	AL1	CKS	OFF									
	16b	8b	8b	8b	16b	16b	8b	8b	16b	16b									

Figure 6-4. Disk Parameter Block Format

where each is a byte or word value, as shown by the 8b or 16b indicator below the field.

The following field abbreviations are used in Figure 6-4:

- o SPT is the total number of sectors per track.
- o BSH is the data allocation block shift factor, determined by the data block allocation size.
- o BLM is the data allocation block mask ( $(2[\text{BSH}-1])$ ).
- o EXM is the extent mask, determined by the data block allocation size and the number of disk blocks.
- o DSM determines the total storage capacity of the disk drive.
- o DRM determines the total number of directory entries that can be in the directory.
- o AL0 and AL1 are the address of the first and last reserved tracks.
- o CKS is the size of the directory check vector.
- o OFF is the number of reserved tracks at the beginning of the (logical) disk.

The values of BSH and BLM determine the data allocation size BLS, which is not an entry in the DPB. Given that the designer has selected a value for BLS, the values of BSH and BLM are shown Table 6-8.

Table 6-8. BSH and BLM Values

				BLS					BSH					BLM					
				1024					3					7					
				2048					4					15					
				4096					5					31					
				8192					6					63					
				16,384					7					127					

where all values are in decimal. The value of EXM depends upon both the BLS and whether the DSM value is less than 256 or greater than 255, as shown in Table 6-9.



## EXM values

DSM>255
---------

6	-	2	7
---	---	---	---



Thus, if  $DRM = 127$  (128 directory entries) and  $BLK = 1024$ , there are 32 directory entries per block, requiring 4 reserved blocks. In this case, the 4 high-order bits of  $AL0$  are set, resulting in the values  $AL0 = 0F0H$  and  $AL1 = 00H$ .

The  $CKS$  value is determined as follows: if the disk drive media is removable, then  $CKS = (DRM+1)/4$ , where  $DRM$  is the last directory entry number. If the media are fixed, then set  $CKS = 0$  (no directory records are checked in this case).

Finally, the  $OFF$  field determines the number of tracks that are skipped at the beginning of the physical disk. This value is automatically added whenever  $SETTRK$  is called and can be used as a mechanism for skipping reserved operating system tracks or for partitioning a large disk into smaller segmented sections.

To complete the discussion of the  $DPB$ , several  $DPHs$  can address the same  $DPB$  if their drive characteristics are identical. Further, the  $DPB$  can be dynamically changed when a new drive is addressed by simply changing the pointer in the  $DPH$ ; because the  $BDOS$  copies the  $DPB$  values to a local area whenever the  $SELDISK$  function is invoked.

Returning back to  $DPH$  for a particular drive, the two address values  $CSV$  and  $ALV$  remain. Both addresses reference an area of uninitialized memory following the  $BIOS$ . The areas must be unique for each drive, and the size of each area is determined by the

The size of the area addressed by  $CSV$  is  $CKS$  bytes, which is sufficient to hold the directory check information for this particular drive. If  $CKS = (DRM+1)/4$ , you must reserve  $(DRM+1)/4$  bytes for directory check use. If  $CKS = 0$ , no storage is reserved.

The size of the area addressed by  $ALV$  is determined by the maximum number of data blocks allowed for this particular disk and is computed as  $(DSM/8)+1$ .

The  $CBIOS$  shown in Appendix B demonstrates an instance of these tables for standard 8-inch, single-density drives. It might be useful to examine this program and compare the tabular values with the definitions given above.

## 6.11 The DISKDEF Macro Library

A macro library called  $DISKDEF$  (shown in Appendix F), greatly simplifies the table construction process. You must have access to the  $MAC$  macro assembler, of course, to use the  $DISKDEF$  facility, while the macro library is included with all CP/M 2 distribution disks.



A BIOS disk definition consists of the following sequence of macro statements:

```
MACLIB      DISKDEF
*****
DISKS      n
DISKDEF    0,...
DISKDEF    1,...
*****
DISKDEF    n-1
*****
ENDEF
```

where the MACLIB statement loads the DISKDEF.LIB file, on the same disk as the BIOS, into MAC's internal tables. The DISKS macro call follows, which specifies the number of drives to be configured with the user's system, where n is an integer in the range 1 to 16. A series of DISKDEF macro calls then follow that define the characteristics of each logical disk, 0 through n-1, corresponding to logical drives A through P. The DISKS and DISKDEF macros generate the in-line fixed data tables described in the previous section and thus must be placed in a nonexecutable portion of the BIOS, typically directly following the BIOS jump vector.

The remaining portion of the BIOS is defined following the DISKDEF macros, with the ENDEF macro call immediately preceding the necessary uninitialized RAM areas that are located in memory above the BIOS.

The DISKDEF macro call takes the form:

```
DISKDEF  dn,fsc,lsc,[skf],bls dks,dir,cks,ofs,[0]
```

where

- o dn is the logical disk number, 0 to n-1.
- o fsc is the first physical sector number (0 or 1).
- o lsc is the last sector number.
- o skf is the optional sector skew factor.
- o bls is the data allocation block size.
- o dks is the number of blocks on the disk.
- o dir is the number of directory entries.
- o cks is the number of checked directory entries.
- o ofs is the track offset to logical track 00.
- o [0] is an optional 1.4 compatibility flag.

The value dn is the drive number being defined with this DISKDEF macro invocation. The fsc parameter accounts for differing sector numbering systems and is usually 0 to 1. The lsc is the last numbered sector on a track. When present, the skf parameter defines the sector skew factor, which is used to create a sector translation table according to the skew.







CP/M Operating System Manual 6.11 The DISKDEF Macro Library

```
DISKDEF 2,0
DISKDEF 3,0
****
ENDEF
```

with all disks having the same parameter values of 26 sectors per track, numbered 1 through 26, with 6 sectors skipped between each access, 1024 bytes per data block, 243 data blocks for a total of 243K-byte disk capacity, 64 checked directory entries, and two operating system tracks.

The DISKS macro generates n DPHs, starting at the DPH table address DPBASE generated by the macro. Each disk header block contains sixteen bytes, as described above, and correspond one-for-one to each of the defined drives. In the four-drive standard system, for example, the DISKS macro generates a table of the form:

```
DPBASE EQU$
DPE0: DW XLTO,0000H,0000H,0000H,DIRBUF,DPB0,CSV0,ALV0
DPE1: DW XLTO,0000H,0000H,0000H,DIRBUF,DPB0,CSV1,ALV1
DPE2: DW XLTO,0000H,0000H,0000H,DIRBUF,DPB0,CSV2,ALV2
DPE3: DW XLTO,0000H,0000H,0000H,DIRBUF,DPB0,CSV3,ALV3
```

where the DPH labels are included for reference purposes to show the beginning table addresses for each drive 0 through 3. The values contained within the DPH are described in detail in the previous section. The check and allocation vector addresses are generated by

Note that if the skf (skew factor) parameter is omitted, or equal to 0, the translation table is omitted and a 0000H value is inserted in the XLTO position of the DPH for the disk. In a subsequent call to perform the logical-to-physical translation, SECTRAN receives a translation table address of DE = 0000H and simply returns the original logical sector from BC in the HL register pair.

A translate table is constructed when the skf parameter is present, and the (nonzero) table address is placed into the corresponding DPHs. The following, for example, is constructed when the standard skew factor skf = 6 is specified in the DISKDEF macro call:

```
XLTO: DB 1,7,13,19,25,5,11,17,23,3,9,15,21
      DB 2,8,14,20,26,6,12,18,24,4,10,16,22
```

Following the ENDEF macro call, a number of uninitialized data areas are defined. These data areas need not be a part of the BIOS that is loaded upon cold start, but must be available between the BIOS and the end of memory. The size of the uninitialized RAM area is determined by EQU statements generated by the ENDEF macro. For a standard four-drive system, the ENDEF macro might produce the following EQU statement:



CP/M Operating System Manual 6.11 The DISKDEF Macro Library

```
4C72 = BEGDAT EQU $
      (data areas)
```

[illegible]

013C	=	DATSIZ	FOU	S-BEGDAT
------	---	--------	-----	----------

which indicates that uninitialized RAM begins at location 4C72H, ends at 4DB0H-1, and occupies 013CH bytes. You must ensure that these addresses are free for use after the system is loaded.

After modification, you can use the STAT program to check drive characteristics, because STAT uses the disk parameter block to decode the drive information. A STAT command of the form:

[illegible]

```
decodes the disk parameter block for drive d ((d=A,...,P)) and
displays the following values:
```

r:	128-byte record capacity
k:	kilobyte drive capacity
d:	32-byte directory entries
c:	checked directory entries
e:	records/extent
b:	records/block

t:	reserved tracks
----	-----------------

Three examples of DISKDEF macro invocations are shown below with corresponding STAT parameter values. The last example produces a full 8-megabyte system.

```
DISKDEF 0,1,58,,2048,256,128,128,2
r=4096, k=512, d=128, c=128, e=256, b=16, s=58, t=2
```

[illegible][illegible]

## 6.12 Sector Blocking and Deblocking

Upon each call to BIOS WRITE entry point, the CP/M BDOS includes information that allows effective sector blocking and deblocking where the host disk subsystem has a sector size that is a multiple of the basic 128-byte unit. The purpose here is to present a general-purpose algorithm that can be included within the BIOS and that uses the BDOS information to perform the operations automatically.



On each call to WRITE, the BDOS provides the following information in register C:

0	=	(normal sector write)
1	=	(write to directory sector)
2	=	(write to the first sector of a new data block)

Condition 0 occurs whenever the next write operation is into a previously written area, such as a random mode record update; when the write is to other than the first sector of an unallocated block; or when the write is not into the directory area. Condition 1 occurs when a write into the directory area is performed. Condition 2 occurs when the first record (only) of a newly allocated data block is written. In most cases, application programs read or write multiple 128-byte sectors in sequence; thus, there is little overhead involved in either operation when blocking and deblocking records, because preread operations can be avoided when writing records.

Appendix G lists the blocking and deblocking algorithms in skeletal form; this file is included on your CP/M disk. Generally, the algorithms map all CP/M sector read operations onto the host disk through an intermediate buffer that is the size of the host disk sector. Throughout the program, values and variables that relate to the CP/M sector involved in a seek operation are prefixed by sek, while those related to the host disk system are prefixed by

define the mapping between CP/M and the host system, and must be changed if other than the sample host system is involved.

The entry points BOOT and WBOOT must contain the initialization code starting on line 57, while the SETDSK entry point must be augmented by the code starting on line 65. Note that although the SETDSK entry point computes and returns the Disk Parameter Header address, it does not physically select the host disk at this point (it is selected later at READHST or WRITEHST). Further, SETTRK, SETTRK, and SETTMA simply store the values, but do not take any other action at this point. SECTRAN performs a trivial function of returning the physical sector number.

The principal entry points are READ and WRITE, starting on lines 110 and 125, respectively. These subroutines take the place of your previous READ and WRITE operations.

The actual physical read or write takes place at either WRITEHST or READHST, where all values have been prepared: hstdsk is the host disk number, hsttrk is the host track number, and hstsec is the host sector number, which may require translation to physical sector number. You must insert code at this point that performs the full sector read or write into or out of the buffer at hstbuf of length hstsiz. All other mapping functions are performed by the algorithms.



[illegible]



Appendix A

The Microcomputer Development System Basic Input/Output System (BIOS)

```
1 ; mds-800 i/o drivers for cp/m 2.2
2 ; (four drive single density version)
3 ;
4 ; version 2.2 february, 1980
5 ;
6 0016 = vers equ 22 ;version 2.2
7 ;
8 ; copyright (c) 1980
9 ; digital research
10 ; box 579, pacific grove
11 ; california, 93950
12 ;
13 ;
14 ffff = true equ 0ffffh ;value of "true"
15 0000 = false equ not true ;"false"
16 0000 = test equ false ;true if test bios
17 ;
18 ; if test
19 bias equ 03400h ;base of ccp in test system
20 endif
21 ; if not test
22 0000 = bias equ 0000h ;generate relocatable cp/m system
23 endif
24 ;
25 1600 = patch equ 1600h
26 ;
27 1600 org patch
28 0000 = cpmb equ $-patch ;base of cpm console processor
29 0806 = bdos equ 806h+cpmb ;basic dos (resident portion)
30 1600 = cpm1 equ $-cpmb ;length (in bytes) of cpm system
31 002c = nsects equ cpm1/128 ;number of sectors to load
32 0002 = offset equ 2 ;number of disk tracks used by cp/m
33 0004 = cdisk equ 0004h ;address of last logged disk on warm start
34 0080 = buff equ 0080h ;default buffer address
```











[illegible]



```

155      ;
156      ; end of controller--independent code, the remaining subroutines
157      ; are tailored to the particular operating environment, and must
158      ; be altered for any system which differs from the intel mds.
159      ;
160      ; the following code assumes the mds monitor exists at 0f800h
161      ; and uses the i/o subroutines within the monitor
162      ;
163      ; we also assume the mds system has four disk drives
164      00fd = revrt    equ    0fdh        ;interrupt revert port
165      00fc = intc     equ    0fch        ;interrupt mask port
166      00f3 = icon     equ    0f3h        ;interrupt control port
167      007E = inte     equ    0111$11110b ;enable rst 0 (warm boot), rst 7 (monitor)
168      ;
169      ; mds monitor equates
170      f800 = mon80    equ    0f800h      ;mds monitor
171      ff0f = rmon80   equ    0ff0fh      ;restart mon80 (boot error)
172      f803 = ci       equ    0f803h      ;console character to reg-a
173      f806 = ri       equ    0f806h      ;reader in to reg-a
174      f809 = co       equ    0f809h      ;console char from c to console out
175      f80c = po       equ    0f80ch      ;punch char from c to punch device
176      f80f = lo       equ    0f80fh      ;list from c to list device
177      f812 = csts     equ    0f812h      ;console status 00/ff to register a
178      ;
179      ; disk ports and commands
180      0078 = base     equ    78h         ;base of disk command io ports
181      0078 = dstat    equ    base       ;disk status (input)
182      0079 = rtype    equ    base+1     ;result type (input)
183      007b = rbyte    equ    base+3     ;result byte (input)
184      ;
185      0079 = ilow     equ    base+1     ;iopb low address (output)
186      007a = ihigh    equ    base+2     ;iopb high address (output)
187      ;
188      0004 = readf     equ    4h         ;read function
189      0006 = writf     equ    6h         ;write function
190      0003 = recal     equ    3h         ;recalibrate drive
191      0004 = iordy     equ    4h         ;i/o finished mask
192      000d = cr        equ    0dh        ;carriage return
193      000a = lf        equ    0ah        ;line-feed
194      ;

```















[illegible]



[illegible]







[illegible]



[illegible]



[illegible]



[illegible]



[illegible]



[illegible]



[illegible]



[illegible]







[illegible]























[illegible]



[illegible]















[illegible]



[illegible]











[illegible]



[illegible]



[illegible]



[illegible]



[illegible]



CP/M 2.0 disk re-definition library

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# Digital Research

Box	579
-----	-----

Pacific Grove, CA

9	3	9	5	0
---	---	---	---	---

CP/M logical disk drives are defined using the

macros given below, where the sequence of calls

i	s	:
---	---	---

disks n

~~diskdef.parameter-list-0~~

0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0
diskdef parameter-list-1																				

[illegible]

```
diskdef parameter-list-n
```

```
endef
```

where  $n$  is the number of logical disk drives attached

to the CP/M system, and parameter-list-i defines the

characteristics of the  $i$ th drive ( $i=0,1,\dots,n-1$ )

\_\_\_\_\_

each parameter-list-*i* takes the form

```
dn,fsc,lsc,[skf],bbs,dks,dir,cks,ofs,[0]
```

where

$dn$  is the disk number  $0, 1, \dots, n-1$

fsc is the first sector number (usually 0 or 1)

```
lsc      is the last sector number on a track
```

```
skf      is optional "skew factor" for sector translate
```

```
bls is the data block size (1024, 2048, ..., 16384)
```

dk is the disk size in bls increments (word)

```
dir    is the number of directory elements (word)
```

```
cks      is the number of dir elements to checksum
```

ofs is the number of tracks to skip (word)















[illegible]



```

196:                                else
197:;;                                generate the translate table
198:nxtsec                          set      0                                ;;next sector to fill
199:nxtbas                          set      0                                ;;moves by one on overflow
200:                                gcd      %sectors,skf
201:;;                                gcdn = gcd(sectors,skew)
202:neltst                          set      sectors/gcdn
203:;;                                neltst is number of elements to generate
204:;;                                before we overlap previous elements
205:nelts                            set      neltst                                ;;counter
206:xlt&dn                        equ      $                                ;;translate table
207:                                rept      sectors                                ;;once for each sector
208:                                if      sectors<256
209:                                ddb      %nxtsec+(fsc)
210:                                else
211:                                ddw      %nxtsec+(fsc)
212:                                endif
213:nxtsec                          set      nxtsec+(skf)
214:                                if      nxtsec>=sectors
215:nxtsec                          set      nxtsec-sectors
216:                                endif
217:nelts                            set      nelts-1
218:                                if      nelts = 0
219:nxtbas                          set      nxtbas+1
220:nxtsec                          set      nxtbas
221:nelts                            set      neltst
222:                                endif
223:                                endm
224:                                endif      ;;end of nul fac test
225:                                endif      ;;end of nul bls test
226:                                endm
227:;;
228:defds                            macro      lab,space
229:lab:                            ds      space
230:                                endm
231:;;
232:lds                            macro      lb,dn,val
233:                                defds      lb&dn,%val&dn
234:                                endm
235:;;

```



[illegible]



Appendix G

## Blocking and Deblocking Algorithms

sector deblocking algorithms for cp/m 2.0

```
utility macro to compute sector mask
```

[illegible]

```
;; compute log2(hb1k), return @x as result
```

```
;; (2 ** @x = hblk on return)
```

```
@v          set          hblk
```

[illegible]

```
;; count right shifts of @y until @x==1
```

7	8	9	0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7
											rept									8

```
if @y = 1
```

```

exitm

```

```
endif
```

```
;; @y is not 1, shift right one position
```

```
@y      set      @y shr 1
```

@x						set				@x + 1
----	--	--	--	--	--	-----	--	--	--	--------

```

endm

```

```

endm

```

$$j$$
[illegible]
$$j$$

```

;      cp/m to host disk constants

```

$$j$$
$$j$$

```
blksiz      equ      2048                                ;cp/m allocation size
```

```
hstsiz      equ      512      ;host disk sector size
```

```
hstspt      equ      20                                ;host disk sectors/trk
```

```
hstblk      equ      hstsiz/128                                ;cp/m sects/host buff
```

```
cpmspt      equ      hstblk * hstspt      ;cp/m sectors/track
```

```
secmsk      equ      hstblk-1      ;sector mask
```

G-1







[illegible]



















[illegible]











[illegible]



[illegible]



## Appendix H

### Glossary

**address:** Number representing the location of a byte in memory. Within CP/M there are two kinds of addresses: logical and physical. A physical address refers to an absolute and unique location within the computer's memory space. A logical address refers to the offset or displacement of a byte in relation to a base location. A standard CP/M program is loaded at address 0100H, the base value; the first instruction of a program has a physical address of 0100H and a relative address or offset of 0H.

**allocation vector (ALV):** An allocation vector is maintained in the BIOS for each logged-in disk drive. A vector consists of a string of bits, one for each block on the drive. The bit corresponding to a particular block is set to one when the block has been allocated and to zero otherwise. The first two bytes of this vector are initialized with the bytes AL0 and AL1 on, thus allocating the directory blocks. CP/M Function 27 returns the allocation vector address.

**AL0, AL1:** Two bytes in the disk parameter block that reserve data blocks for the directory. These two bytes are copied into the first two bytes of the allocation vector when a drive is logged in. See

**ALV:** See **allocation vector**.

**ambiguous filename:** Filename that contains either of the CP/M wildcard characters, ? or \*, in the primary filename, filetype, or both. When you replace characters in a filename with these wildcard characters, you create an ambiguous filename and can easily reference more than one CP/M file in a single command line.

**American Standard Code for Information Interchange:** See **ASCII**.

**applications program:** Program designed to solve a specific problem. Typical applications programs are business accounting packages, word processing (editing) programs and mailing list programs.

**archive attribute:** File attribute controlled by the high-order bit of the t3 byte (FCB+11) in a directory element. This attribute is set if the file has been archived.

**argument:** Symbol, usually a letter, indicating a place into which you can substitute a number, letter, or name to give an appropriate meaning to the formula in question.

**ASCII:** American Standard Code for Information Interchange. ASCII is a standard set of seven-bit numeric character codes used to represent characters in memory. Each character requires one byte of memory with the high-order bit usually set to zero. Characters can



be numbers, letters, and symbols. An ASCII file can be intelligibly displayed on the video screen or printed on paper.

**assembler:** Program that translates assembly language into the binary machine code. Assembly language is simply a set of mnemonics used to designate the instruction set of the CPU. See **ASM** in Section 3 of this manual.

**back-up:** Copy of a disk or file made for safekeeping, or the creation of the duplicate disk or file.

**Basic Disk Operating System:** See **BDOS**.

**BDOS:** Basic Disk Operating System. The BDOS module of the CP/M operating system provides an interface for a user program to the operating system. This interface is in the form of a set of function calls which may be made to the BDOS through calls to location 0005H in page zero. The user program specifies the number of the desired function in register C. User programs running under CP/M should use BDOS functions for all I/O operations to remain compatible with other CP/M systems and future releases. The BDOS normally resides in high memory directly below the BIOS.

**bias:** Address value which when added to the origin address of your BIOS module produces 1F80H, the address of the BIOS module in the MOVCPM image. There is also a bias value that when added to the BOOT module origin produces 0900H, the address of the BOOT module in the MOVCPM image. These values are used to patch the BIOS and BOOT modules under DDT or STD when you patch a CP/M system. If you do not, the patched system may fail to function.

**binary:** Base 2 numbering system. A binary digit can have one of two values: 0 or 1. Binary numbers are used in computers because the hardware can most easily exhibit two states: off and on. Generally, a bit in memory represents one binary digit.

**Basic Input/Output System:** See **BIOS**.

**BIOS:** Basic Input/Output System. The BIOS is the only hardware-dependent module of the CP/M system. It provides the BDOS with a set of primitive I/O operations. The BIOS is an assembly language module usually written by the user, hardware manufacturer, or independent software vendor, and is the key to CP/M's portability. The BIOS interfaces the CP/M system to its hardware environment through a standardized jump table at the front of the BIOS routine and through a set of disk parameter tables which define the disk environment. Thus, the BIOS provides CP/M with a completely table-driven I/O system.

**BIOS base:** Lowest address of the BIOS module in memory, that by definition must be the first entry point in the BIOS jump table.



**bit:** Switch in memory that can be set to on (1) or off (0). Bits are grouped into bytes, eight bits to a byte, which is the smallest directly addressable unit in an Intel 8080 or Zilog Z80. By common convention, the bits in a byte are numbered from right, 0 for the low-order bit, to left, 7 for the high-order bit. Bit values are often represented in hexadecimal notation by grouping the bits from the low-order bit in groups of four. Each group of four bits can have a value from 0 to 15 and thus can easily be represented by one hexadecimal digit.

**BLM:** See **block mask**.

**block:** Basic unit of disk space allocation. Each disk drive has a fixed block size (BLS) defined in its disk parameter block in the BIOS. A block can consist of 1K, 2K, 4K, 8K, or 16K consecutive bytes. Blocks are numbered relative to zero so that each block is unique and has a byte displacement in a file equal to the block number times the block size.

**block mask (BLM):** Byte value in the disk parameter block at DPB + 3. The block mask is always one less than the number of 128 byte sectors that are in one block. Note that  $BLM = (2 ** BSH) - 1$ .

**block shift (BSH):** Byte parameter in the disk parameter block at DPB + 2. Block shift and block mask (BLM) values are determined by the block size (BLS). Note that  $BLM = (2 ** BSH) - 1$ .

**blocking and deblocking algorithm:** If the host sector size is larger than 128 bytes, usually 256, 512, 1024, or 2048 bytes. When the host sector size is larger than 128 bytes, host sectors must be buffered in memory and the 128-byte CP/M sectors must be blocked and deblocked by adding an additional module, the blocking and deblocking algorithm, between the BIOS disk I/O routines and the actual disk I/O. The host sector size must be an even multiple of 128 bytes for the algorithm to work correctly. The blocking and deblocking algorithm allows the BDOS and BIOS to function exactly as if the entire disk consisted only of 128-byte sectors, as in the standard CP/M installation.

**BLS:** Block size in bytes. See **block**.

**boot:** Process of loading an operating system into memory. A boot program is a small piece of code that is automatically executed when you power-up or reset your computer. The boot program loads the rest of the operating system into memory in a manner similar to a person pulling himself up by his own bootstraps. This process is sometimes called a cold boot or cold start. Bootstrap procedures vary from system to system. The boot program must be customized for the memory size and hardware environment that the operating system manages. Typically, the boot resides on the first sector of the system tracks on your system disk. When executed, the boot loads the remaining sectors of the system tracks into high memory at the location for which the CP/M system has been configured. Finally, the boot transfers execution to the boot entry point in the BIOS jump table so that the system can initialize itself. In this case,



the boot program should be placed at 900H in the SYSGEN image. Alternatively, the boot program may be located in ROM.

**bootstrap:** See **boot**.

**BSH:** See **block shift**.

**BTREE:** General purpose file access method that has become the standard organization for indexes in large data base systems. BTREE provides near optimum performance over the full range of file operations, such as insertion, deletion, search, and search next.

**buffer:** Area of memory that temporarily stores data during the transfer of information.

**built-in commands:** Commands that permanently reside in memory. They respond quickly because they are not accessed from a disk.

**byte:** Unit of memory or disk storage containing eight bits. A byte can represent a binary number between 0 and 255, and is the smallest unit of memory that can be addressed directly in 8-bit CPUs such as the Intel 8080 or Zilog Z80.

**CCP:** Console Command Processor. The CCP is a module of the CP/M operating system. It is loaded directly below the BDOS module and interprets and executes commands typed by the console user. Usually these commands are programs that the CCP loads and calls. Upon

has not overwritten it. If it has, the program can reload the CCP into memory by a warm boot operation initiated by either a jump to zero, BDOS system reset (Function 0), or a cold boot. Except for its location in high memory, the CCP works like any other standard CP/M program; that is, it makes only BDOS function calls for its I/O operations.

**CCP base:** Lowest address of the CCP module in memory. This term sometimes refers to the base of the CP/M system in memory, as the CCP is normally the lowest CP/M module in high memory.

**checksum vector (CSV):** Contiguous data area in the BIOS, with one byte for each directory sector to be checked, that is, CKS bytes. See **CKS**. A checksum vector is initialized and maintained for each logged-in drive. Each directory access by the system results in a checksum calculation that is compared with the one in the checksum vector. If there is a discrepancy, the drive is set to Read-Only status. This feature prevents the user from inadvertently switching disks without logging in the new disk. If the new disk is not logged-in, it is treated the same as the old one, and data on it might be destroyed if writing is done.

**CKS:** Number of directory records to be checked summed on directory accesses. This is a parameter in the disk parameter block located in the BIOS. If the value of CKS is zero, then no directory records are checked. CKS is also a parameter in the diskdef macro library, where it is the actual number of directory elements to be checked rather than the number of directory records.



CP/M Operating System Manual H Glossary

**cold boot:** See **boot**. Cold boot also refers to a jump to the boot entry point in the BIOS jump table.

**COM:** Filetype for a CP/M command file. See **command file**.

**command:** CP/M command line. In general, a CP/M command line has three parts: the command keyword, command tail, and a carriage return. To execute a command, enter a CP/M command line directly after the CP/M prompt at the console and press the carriage return or enter key.

**command file:** Executable program file of filetype COM. A command file is a machine language object module ready to be loaded and executed at the absolute address of 0100H. To execute a command file, enter its primary filename as the command keyword in a CP/M command line.

**command keyword:** Name that identifies a CP/M command, usually the primary filename of a file of type COM, or a built-in command. The command keyword precedes the command tail and the carriage return in the command line.

**command syntax:** Statement that defines the correct way to enter a command. The correct structure generally includes the command keyword, the command tail, and a carriage return. A syntax line usually contains symbols that you should replace with actual values

**command tail:** Part of a command that follows the command keyword in the command line. The command tail can include a drive specification, a filename and filetype, and options or parameters. Some commands do not require a command tail.

**CON:** Mnemonic that represents the CP/M console device. For example, the CP/M command PIP CON:=TEST.SUB displays the file TEST.SUB on the console device. The explanation of the STAT command tells how to assign the logical device CON: to various physical devices. See **console**.

**concatenate:** Name of the PIP operation that copies two or more separate files into one new file in the the specified sequence.

**concurrency:** Execution of two processes or operations simultaneously.

**CONIN:** BIOS entry point to a routine that reads a character from the console device.

**CONOUT:** BIOS entry point to a routine that sends a character to the console device.



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-----



**data file:** File containing information that will be processed by a program.

**deblocking:** See **blocking & deblocking algorithm**.

**default:** Currently selected disk drive and user number. Any command that does not specify a disk drive or a user number references the default disk drive and user number. When CP/M is first invoked, the default disk drive is drive A, and the default user number is 0.

**default buffer:** Default 128-byte buffer maintained at 0080H in page zero. When the CCP loads a COM file, this buffer is initialized to the command tail; that is, any characters typed after the COM file name are loaded into the buffer. The first byte at 0080H contains the length of the command tail, while the command tail itself begins at 0081H. The command tail is terminated by a byte containing a binary zero value. The I command under DDT and SID initializes this buffer in the same way as the CCP.

**default FCB:** Two default FCBs are maintained by the CCP at 005CH and 006CH in page zero. The first default FCB is initialized from the first delimited field in the command tail. The second default FCB is initialized from the next field in the command tail.

**delimiter:** Special characters that separate different items in a filename. The CCP recognizes the following characters as delimiters: . : = ; < > \_ , blank, and carriage return. Several CP/M commands also treat the following as delimiter characters: [ ] ( ) \$ . It is advisable to avoid the use of delimiter characters and lower-case characters in CP/M filenames.

**DIR:** Parameter in the diskdef macro library that specifies the number of directory elements on the drive.

**DIR attribute:** File attribute. A file with the DIR attribute can be displayed by a DIR command. The file can be accessed from the default user number and drive only.

**DIRBUF:** 128-byte scratchpad area for directory operations, usually located at the end of the BIOS. DIRBUF is used by the BDOS during its directory operations. DIRBUF also refers to the two-byte address of this scratchpad buffer in the disk parameter header at DPbase + 8 bytes.

**directory:** Portion of a disk that contains entries for each file on the disk. In response to the DIR command, CP/M displays the filenames stored in the directory. The directory also contains the locations of the blocks allocated to the files. Each file directory element is in the form of a 32-byte FCB, although one file can have several elements, depending on its size. The maximum number of directory elements supported is specified by the drive's disk parameter block value for DRM.



**directory element:** Data structure. Each file on a disk has one or more 32-byte directory elements associated with it. There are four directory elements per directory sector. Directory elements can also be referred to as directory FCBs.

**directory entry:** File entry displayed by the DIR command. Sometimes this term refers to a physical directory element.

**disk, diskette:** Magnetic media used for mass storage in a computer system. Programs and data are recorded on the disk in the same way music can be recorded on cassette tape. The CP/M operating system must be initially loaded from disk when the computer is turned on. Diskette refers to smaller capacity removable floppy diskettes, while disk may refer to either a diskette, removable cartridge disk, or fixed hard disk. Hard disk capacities range from five to several hundred megabytes of storage.

**diskdef macro library:** Library of code that when used with MAC, the Digital Research macro assembler, creates disk definition tables such as the DPB and DPH automatically.

**disk drive:** Peripheral device that reads and writes information on disk. CP/M assigns a letter to each drive under its control. For example, CP/M may refer to the drives in a four-drive system as A, B, C, and D.

**disk parameter block (DPB):** Data structure referenced by one or more disk parameter headers. Characteristics in the fields listed below:

- SPT is the total number of sectors per track.
- BSH is the data allocation block shift factor.
- BLM is the data allocation block mask.
- EXM is the extent mask determined by BLS and DSM.
- DSM is the maximum data block number.
- DRM is the maximum number of directory entries--1.
- AL0 reserves directory blocks.
- AL1 reserves directory blocks.
- CKS is the number of directory sectors check summed.
- OFF is the number of reserved system tracks.

The address of the disk parameter block is located in the disk parameter header at DPbase +0AH. CP/M Function 31 returns the DPB address. Drives with the same characteristics can use the same disk parameter header, and thus the same DPB. However, drives with different characteristics must each have their own disk parameter header and disk parameter blocks. When the BDOS calls the SEIDSK entry point in the BIOS, SEIDSK must return the address of the drive's disk parameter header in register HL.

**disk parameter header (DPH):** Data structure that contains information about the disk drive and provides a scratchpad area for certain BDOS operations. The disk parameter header contains six bytes of scratchpad area for the BDOS, and the following five 2-byte parameters:



Given  $n$  disk drives, the disk parameter headers are arranged in a table whose first row of 16 bytes corresponds to drive 0, with the last row corresponding to drive  $n-1$ .

**DMA:** Direct Memory Access. DMA is a method of transferring data from the disk into memory directly. In a CP/M system, the BDOS calls the BIOS entry point READ to read a sector from the disk into the currently selected DMA address. The DMA address must be the address of a 128-byte buffer in memory, either the default buffer at 0080H in page zero, or a user-assigned buffer in the TPA. Similarly, the BDOS calls the BIOS entry point WRITE to write the record at the current DMA address to the disk.

```
DPH: See disk parameter header.
```

**DSM:** 2-byte parameter of the disk parameter block at DPB + 5. DSM is the maximum data block number supported by the drive. The product BLS times (DSM+1) is the total number of bytes held by the drive. This must not exceed the capacity of the physical disk less the reserved system tracks.

**EX:** Extent number field in an FCB. See **extent**.

**execute a program:** Start the processing of executable code.



**EXM:** See **extent mask**.

**extent:** 16K consecutive bytes in a file. Extents are numbered from 0 to 31. One extent can contain 1, 2, 4, 8, or 16 blocks. EX is the extent number field of an FCB and is a one-byte field at FCB + 12, where FCB labels the first byte in the FCB. Depending on the block size (BLS) and the maximum data block number (DSM), an FCB can contain 1, 2, 4, 8, or 16 extents. The EX field is normally set to 0 by the user but contains the current extent number during file I/O. The term FCB folding describes FCBs containing more than one extent. In CP/M version 1.4, each FCB contained only one extent. Users attempting to perform random record I/O and maintain CP/M 1.4 compatibility should be aware of the implications of this difference. See **CP/M 1.4 compatibility**.

**extent mask (EXM):** A byte parameter in the disk parameter block located at DPB + 3. The value of EXM is determined by the block size (BLS) and whether the maximum data block number (DSM) exceeds 255. There are EXM + 1 extents per directory FCB.

**FCB:** See **File Control Block**.

**file:** Collection of characters, instructions, or data that can be referenced by a unique identifier. Files are usually stored on various types of media, such as disk, or magnetic tape. A CP/M file is identified by a file specification and resides on disk as a file and can contain either binary or ASCII data. Binary files contain bytes of data that can vary in value from 0H to 0FFH. ASCII files contain sequences of character codes delineated by a carriage return and line-feed combination; normally byte values range from 0H to 7FH. The directory maps the file as a series of physical blocks. Although files are defined as a sequence of consecutive logical records, these records can not reside in consecutive sectors on the disk. See also **block**, **directory**, **extent**, **record**, and **sector**.

**File Control Block (FCB):** Structure used for accessing files on disk. Contains the drive, filename, filetype, and other information describing a file to be accessed or created on the disk. A file control block consists of 36 consecutive bytes specified by the user for file I/O functions. FCB can also refer to a directory element in the directory portion of the allocated disk space. These contain the same first 32 bytes of the FCB, but lack the current record and random record number bytes.

**filename:** Name assigned to a file. A filename can include a primary filename of one to eight characters; a filetype of zero to three characters. A period separates the primary filename from the filetype.

**file specification:** Unique file identifier. A complete CP/M file specification includes a disk drive specification followed by a colon, d:, a primary filename of one to eight characters, a period, and a filetype of zero to three characters. For example, b:example.tex is a complete CP/M file specification.



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same in its actual physical form. For example, a hard disk can occupy one physical drive, yet you can divide the available storage on it to appear to the user as if it were in several different drives. These apparent drives are the logical drives.

**logical sector:** See **sector**.

**logical-to-physical sector translation table:** See **XLT**.

**LSC:** Diskdef macro library parameter specifying the last physical sector number.

**LST:** Logical CP/M list device, usually a printer. The CP/M list device is an output-only device referenced through the LIST and LISTST entry points of the BIOS. The STAT command allows assignment of LST: to one of the physical devices: ITTY:, CRT:, LPT:, or UL1:, provided these devices and the IOBYTE are implemented in the LIST and LISTST entry points of your CP/M BIOS module. The CP/NET command NETWORK allows assignment of LST: to a list device on a network master. For example, PIP LST:=TEST.SUB prints the file TEST.SUB on the list device.

**macro assembler:** Assembler code translator providing macro processing facilities. Macro definitions allow groups of instructions to be stored and substituted in the source program as the macro names are encountered. Definitions and invocations can be nested and macro parameters can be formed to pass arbitrary strings

**megabyte:** Over one million bytes; 1024 kilobytes. See **byte**, and **kilobyte**.

**microprocessor:** Silicon chip that is the central processing unit (CPU) of the microcomputer. The Intel 8080 and the Zilog Z80 are microprocessors commonly used in CP/M systems.

**MOVCPM image:** Memory image of the CP/M system created by MOVCPM. This image can be saved as a disk file using the SAVE command or placed on the system tracks using the SYSGEN command without specifying a source drive. This image varies, depending on the presence of a one-sector or two-sector boot. If the boot is less than 128 bytes (one sector), the boot begins at 0900H, the CP/M system at 0980H, and the BIOS at 1F80H. Otherwise, the boot is at 0900H, the CP/M system at 1000H, and the BIOS at 2000H. In a CP/M 1.4 system with a one-sector boot, the addresses are the same as for the CP/M 2 system--except that the BIOS begins at 1E80H instead of 1F80H.

**MP/M:** Multi-Programming Monitor control program. A microcomputer operating system supporting multi-terminal access with multi-programming at each terminal.

**multi-programming:** The capability of initiating and executing more than one program at a time. These programs, usually called processes, are time-shared, each receiving a slice of CPU time on a round-robin basis. See **concurrency**.



**nibble:** One half of a byte, usually the high-order or low-order 4 bits in a byte.

**OFF:** Two-byte parameter in the disk parameter block at DPB + 13 bytes. This value specifies the number of reserved system tracks. The disk directory begins in the first sector of track OFF.

**OFS:** Diskdef macro library parameter specifying the number of reserved system tracks. See **OFF**.

**operating system:** Collection of programs that supervises the execution of other programs and the management of computer resources. An operating system provides an orderly input/output environment between the computer and its peripheral devices. It enables user-written programs to execute safely. An operating system standardizes the use of computer resources for the programs running under it.

**option:** One of many parameters that can be part of a command tail. Use options to specify additional conditions for a command's execution.

**output:** Data that is sent to the console, disk, or printer.

**page:** 256 consecutive bytes in memory beginning on a page boundary, whose base address is a multiple of 256 (100H) bytes. In hex notation, pages always begin at an address with a least significant

**page relocatable program:** See **PRL**.

**page zero:** Memory region between 0000H and 0100H used to hold critical system parameters. Page zero functions primarily as an interface region between user programs and the CP/M BDOS module. Note that in non-standard systems this region is the base page of the system and represents the first 256 bytes of memory used by the CP/M system and user programs running under it.

**parameter:** Value in the command tail that provides additional information for the command. Technically, a parameter is a required element of a command.

**peripheral devices:** Devices external to the CPU. For example, terminals, printers, and disk drives are common peripheral devices that are not part of the processor but are used in conjunction with it.

**physical:** Characteristic of computer components, generally hardware, that actually exist. In programs, physical components can be represented by logical components.

**primary filename:** First 8 characters of a filename. The primary filename is a unique name that helps the user identify the file contents. A primary filename contains one to eight characters and can include any letter or number and some special characters. The



primary filename follows the optional drive specification and precedes the optional filetype.

**PRL:** Page relocatable program. A page relocatable program is stored on disk with a PRL filetype. Page relocatable programs are easily relocated to any page boundary and thus are suitable for execution in a nonbanked MP/M system.

**program:** Series of coded instructions that performs specific tasks when executed by a computer. A program can be written in a processor-specific language or a high-level language that can be implemented on a number of different processors.

**prompt:** Any characters displayed on the video screen to help the user decide what the next appropriate action is. A system prompt is a special prompt displayed by the operating system. The alphabetic character indicates the default drive. Some applications programs have their own special prompts. See **CP/M prompt**.

**PUN:** Logical CP/M punch device. The punch device is an output-only device accessed through the PUNCH entry point of the BIOS. In certain implementations, PUN: can be a serial device such as a modem.

**PUNCH:** BIOS entry point to a routine that sends a character to the punch device.

**READER:** Logical CP/M reader device. The reader device is an input-only device accessed through the READER entry point in the BIOS. See **PUN:**.

**READ:** Entry point in the BIOS to a routine that reads 128 bytes from the currently selected drive, track, and sector into the current DMA address.

**READER:** Entry point to a routine in the BIOS that reads the next character from the currently assigned reader device.

**Read-Only (R/O):** Attribute that can be assigned to a disk file or a disk drive. When assigned to a file, the Read-Only attribute allows you to read from that file but not write to it. When assigned to a drive, the Read-Only attribute allows you to read any file on the disk, but prevents you from adding a new file, erasing or changing a file, renaming a file, or writing on the disk. The **STAT** command can set a file or a drive to Read-Only. Every file and drive is either Read-Only or Read-Write. The default setting for drives and files is Read-Write, but an error in resetting the disk or changing media automatically sets the drive to Read-Only until the error is corrected. See also **ROM**.

**Read-Write (R/W):** Attribute that can be assigned to a disk file or a disk drive. The Read-Write attribute allows you to read from and write to a specific Read-Write file or to any file on a disk that is in a drive set to Read-Write. A file or drive can be set to either Read-Only or Read-Write.



**record:** Group of bytes in a file. A physical record consists of 128 bytes and is the basic unit of data transfer between the operating system and the application program. A logical record might vary in length and is used to represent a unit of information. Two 64-byte employee records can be stored in one 128-byte physical record. Records are grouped together to form a file.

**recursive procedure:** Code that can call itself during execution.

**reentrant procedure:** Code that can be called by one process while another is already executing it. Thus, reentrant code can be shared between different users. Reentrant procedures must not be self-modifying; that is, they must be pure code and not contain data. The data for reentrant procedures can be kept in a separate data area or placed on the stack.

**restart (RST):** One-byte call instruction usually used during interrupt sequences and for debugger break pointing. There are eight restart locations, RST 0 through RST 7, whose addresses are given by the product of 8 times the restart number.

**R/O:** See **Read-Only**.

**ROM:** Read-Only memory. This memory can be read but not written and so is suitable for code and preinitialized data areas only.

**RST:** See **restart**.

**R/W:** See **Read-Write**.

**sector:** In a CP/M system, a sector is always 128 consecutive bytes. A sector is the basic unit of data read and written on the disk by the BIOS. A sector can be one 128-byte record in a file or a sector of the directory. The BDOS always requests a logical sector number between 0 and (SPT-1). This is typically translated into a physical sector by the BIOS entry point SECTRAN. In some disk subsystems, the disk sector size is larger than 128 bytes, usually a power of two, such as 256, 512, 1024, or 2048 bytes. These disk sectors are always referred to as host sectors in CP/M documentation and should not be confused with other references to sectors, in which cases the CP/M 128-byte sectors should be assumed. When the host sector size is larger than 128 bytes, host sectors must be buffered in memory and the 128-byte CP/M sectors must be blocked and deblocked from them. This can be done by adding an additional module, the blocking and deblocking algorithm, between the BIOS disk I/O routines and the actual disk I/O.

**sectors per track (SPT):** A two-byte parameter in the disk parameter block at DPB + 0. The BDOS makes calls to the BIOS entry point SECTRAN with logical sector numbers ranging between 0 and (SPT - 1) in register BC.

**SECTRAN:** Entry point to a routine in the BIOS that performs logical-to-physical sector translation for the BDOS.



Logical:	0	1	2	3	4	5	6	...	25
----------	---	---	---	---	---	---	---	-----	----

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**spooling:** Process of accumulating printer output in a file while the printer is busy. The file is printed when the printer becomes free; a program does not have to wait for the slow printing process.

**SPT:** See **sectors per track**.

**stack:** Reserved area of memory where the processor saves the return address when a call instruction is received. When a return instruction is encountered, the processor restores the current address on the stack to the program counter. Data such as the contents of the registers can also be saved on the stack. The push instruction places data on the stack and the pop instruction removes it. An item is pushed onto the stack by decrementing the stack pointer (SP) by 2 and writing the item at the SP address. In other words, the stack grows downward in memory.

**syntax:** Format for entering a given command.

**SYS:** See **system attribute**.

**SYSGEN image:** Memory image of the CP/M system created by SYSGEN when a destination drive is not specified. This is the same as the MOVCPM image that can be read by SYSGEN if a source drive is not specified. See **MOVCPM image**.

**system attribute (SYS):** File attribute. You can give a file the system attribute by using the SYS option in the STAT command or by with the SYS attribute is not displayed in response to a DIR command. If you give a file with user number 0 the SYS attribute, you can read and execute that file from any user number on the same drive. Use this feature to make your commonly used programs available under any user number.

**system prompt:** Symbol displayed by the operating system indicating that the system is ready to receive input. See **prompt** and **CP/M prompt**.

**system tracks:** Tracks reserved on the disk for the CP/M system. The number of system tracks is specified by the parameter OFF in the disk parameter block (DPB). The system tracks for a drive always precede its data tracks. The command SYSGEN copies the CP/M system from the system tracks to memory, and vice versa. The standard SYSGEN utility copies 26 sectors from track 0 and 26 sectors from track 1. When the system tracks contain additional sectors or tracks to be copied, a customized SYSGEN must be used.

**terminal:** See **console**.

**TPA:** Transient Program Area. Area in memory where user programs run and store data. This area is a region of memory beginning at 0100H and extending to the base of the CP/M system in high memory. The first module of the CP/M system is the CCP, which can be overwritten by a user program. If so, the TPA is extended to the base of the CP/M BDOS module. If the CCP is overwritten, the user



program must terminate with either a system reset (Function 0) call or a jump to location zero in page zero. The address of the base of the CP/M BDOS is stored in location 0006H in page zero least significant byte first.

**track:** Data on the disk media is accessed by combination of track and sector numbers. Tracks form concentric rings on the disk; the standard IBM single-density disks have 77 tracks. Each track consists of a fixed number of numbered sectors. Tracks are numbered from zero to one less than the number of tracks on the disk.

**Transient Program Area:** See TPA.

**upward compatible:** Term meaning that a program created for the previously released operating system, or compiler, runs under the newly released version of the same operating system.

**USER:** Term used in CP/M and MP/M systems to distinguish distinct regions of the directory.

**user number:** Number assigned to files in the disk directory so that different users need only deal with their own files and have their own directories, even though they are all working from the same disk. In CP/M, files can be divided into 16 user groups.

**utility:** Tool. Program that enables the user to perform certain operations, such as copying files, erasing files, and editing files.

users.

**vector:** Location in memory. An entry point into the operating system used for making system calls or interrupt handling.

**warm start:** Program termination by a jump to the warm start vector at location 0000H, a system reset (BDOS Function 0), or a CTRL-C typed at the keyboard. A warm start reinitializes the disk subsystem and returns control to the CP/M operating system at the CCP level. The warm start vector is simply a jump to the WBOOT entry point in the BIOS.

**WBOOT:** Entry point to a routine in the BIOS used when a warm start occurs. A warm start is performed when a user program branches to location 0000H, when the CPU is reset from the front panel, or when the user types CTRL-C. The CCP and BDOS are reloaded from the system tracks of drive A.

**wildcard characters:** Special characters that match certain specified items. In CP/M there are two wildcard characters: ? and \*. The ? can be substituted for any single character in a filename, and the \* can be substituted for the primary filename, the filetype, or both. By placing wildcard characters in filenames, the user creates an ambiguous filename and can quickly reference one or more files.







## Appendix I

### CP/M Error Messages

Messages come from several different sources. CP/M displays error messages when there are errors in calls to the Basic Disk Operating System (BDOS). CP/M also displays messages when there are errors in command lines. Each utility supplied with CP/M has its own set of messages. The following lists CP/M messages and utility messages. One might see messages other than those listed here if one is running an application program. Check the application program's documentation for explanations of those messages.

Table I-1. CP/M Error Messages

Message	Meaning
---------	---------

?	
---	--

DDT. This message has four possible meanings:

- o DDT does not understand the assembly language instruction.
- o A checksum error occurred in a HEX file.
- o The assembler/disassembler was overlayed.

ABORTED

PIP. You stopped a PIP operation by pressing a key.

#### ASM Error Messages

D Data error: data statement element cannot be placed in specified data area.

E Expression error: expression cannot be evaluated during assembly.

I Label error: label cannot appear in this context (might be duplicate label).



Message	Meaning
---------	---------

N	Not implemented:	unimplemented
	features, such as	macros, are
	trapped.	

```
0  Overflow: expression is too complex
to evaluate.
```

```
P Phase error: label value changes on
two passes through assembly.
```

```
R   Register error:  the value specified
   as a register is incompatible with
   the code.
```

```
S      Syntax error:      improperly formed
      expression.
```

```
U Undefined label: label used does not
  exist.
```

```
Value error: improperly formed
```

BAD DELIMITER

STAT.	Check	command	line	for	typing
errors.					

Bad Load
----------

CCP error message, or SAVE error message.

Bdos Err On d:

Basic Disk Operating System error on the designated drive: CP/M replaces d: with the drive specification of the drive where the error occurred. This message is followed by one of the four phrases in the situations described below.



Message	Meaning
1. The first message is "I am a doctor"	1. The first message is "I am a doctor"
2. The second message is "I am a doctor"	2. The second message is "I am a doctor"
3. The third message is "I am a doctor"	3. The third message is "I am a doctor"
4. The fourth message is "I am a doctor"	4. The fourth message is "I am a doctor"
5. The fifth message is "I am a doctor"	5. The fifth message is "I am a doctor"
6. The sixth message is "I am a doctor"	6. The sixth message is "I am a doctor"
7. The seventh message is "I am a doctor"	7. The seventh message is "I am a doctor"
8. The eighth message is "I am a doctor"	8. The eighth message is "I am a doctor"
9. The ninth message is "I am a doctor"	9. The ninth message is "I am a doctor"
10. The tenth message is "I am a doctor"	10. The tenth message is "I am a doctor"

```
Bdos Err On d: Bad Sector
```

This message appears when CP/M finds no disk in the drive, when the disk is improperly formatted, when the drive latch is open, or when power to the drive is off. Check for one of these situations and try again. This could also indicate a hardware problem or a worn or improperly formatted disk. Press ^C to terminate the program and return to CP/M, or press RETURN to ignore the error.

```
Bdos Err On d: File R/O
```

You tried to erase, rename, or set file attributes on a Read-Only file. The file should first be set to Read-Write (R/W) with the command: `STAT filespec $R/W`.

Bdos	Err	On	d:	R/O
------	-----	----	----	-----

```
Drive has been assigned Read-Only status
with a STAT command, or the disk in the
drive has been changed without being
initialized with a ^C. CP/M terminates
the current program as soon as you press
any key.
```

Bdos Err on d: Select

CP/M received a command line specifying a nonexistent drive. CP/M terminates the current program as soon as you press any key. Press RETURN or CTRL-C to recover.

```
Break "x" at c
```

ED. "x" is one of the symbols described below and c is the command letter being executed when the error occurred.

#	Search failure. ED cannot find the string specified in an F, S, or N command.
---	---











PIP. The destination file specified in a PIP command already exists and it is Read-Only. If you type Y, the destination file is deleted before the file copy is done.

I-6
-----







[illegible]

```
LOAD.  Cannot find source file.  Check
disk directory.
```

LOAD. Caused by an error code returned by a BDOS function call.

LOAD. Destination disk is full.

LOAD. The address of a record was too far from the address of the previously-processed record. This is an internal limitation of LOAD, but it can be 012345678901234567890123456789012345678901234567890  
file into memory, then use a SAVE command to store the memory image file on disk.

LOAD. Disk directory is full.

SUBMIT. The SUBMIT program displays its messages in the format shown above, where nnn represents the line number of the SUBMIT file. Refer to the message following the line number.

```
ED. Disk or directory is full, and ED
cannot write anything more on the disk.
This is a fatal error, so make sure there
is enough space on the disk to hold a
second copy of the file before invoking
ED.
```







[illegible]

```
PIP. An input file that you have
specified does not exist.
```

```
ED.  You typed the ED command without a
filename.  Reenter the ED command followed
by the name of the file you want to edit
or create.
```

DDT. The ?? indicates DDT does not know how to represent the hexadecimal value dd encountered at address hhhh in 8080 assembly language. dd is not an 8080 machine instruction opcode.

```
DDT.  There is not enough memory to load
the file specified in an R or E command.
```

Invalid Assignment

```
STAT.  You specified an invalid drive or
file assignment, or misspelled a device
name.  This error message might be
followed by a list of the valid file
assignments that can follow a filename.
If an invalid drive assignment was
attempted the message Use: d:=RO is
displayed, showing the proper syntax for
drive assignments.
```



[illegible]

SUBMIT. The only valid control characters in the SUBMIT files of the type SUB are ^ A through ^ Z. Note that in a SUBMIT file the control character is represented by typing the circumflex, ^, not by pressing the control key.

P1P.	An invalid HEX digit has been encountered while reading a HEX file. The HEX file with the invalid HEX digit should be corrected, probably by recreating the HEX file.
------	---

STAT. Might appear if you follow the drive specification with anything except

CP/M received a command line specifying a nonexistent drive, or the disk in the drive is improperly formatted. CP/M terminates the current program as soon as you press any key.

```
SYSGEN.  SYSGEN recognizes only drives A,
B, C, and D as valid destinations for
system generation.
```

STAT.	Appears if you do not specify RO, RW, DIR, or SYS.
-------	--



[illegible]

```
PIP. The format of your PIP command is
illegal. See the description of the PIP
command.
```

INVALID HEX DIGIT
LOAD ADDRESS hhhh
ERROR ADDRESS hhhh
BYTES READ:
hhhh

```
LOAD.  File contains incorrect HEX digit.
```

I	N	V	A	L	I	D		M	E	M	O	R	Y		S	I	Z	E
---	---	---	---	---	---	---	--	---	---	---	---	---	---	--	---	---	---	---

MOVCPM. Specify a value less than 64K or your computer's actual memory size.
---

INVALID SEPARATOR

```
PIP. You have placed an invalid character
for a separator between two input
filenames.
```

I	N	V	A	L	I	D		U	S	E	R		N	U	M	B	E	R
---	---	---	---	---	---	---	--	---	---	---	---	--	---	---	---	---	---	---

```
PIP. You have specified a user number
greater than 15. User numbers are in the
range 0 to 15.
```

n?	
----	--

USER. You specified a number greater than fifteen for a user area number. For example, if you type USER 18<cr>, the screen displays 18?.











Message	Meaning
---------	---------

SUBMIT. Within the SUBMIT file of type sub, valid parameters are \$0 through \$9.

PARAMETER ERROR, TYPE RETURN TO IGNORE

SYSGEN.	If you press RETURN, SYSGEN
proceeds without	processing the invalid
parameter.	

QUIT NOT FOUND

PIP. The string argument to a Q parameter was not found in your input file.

Read	error
------	-------

TYPE. An error occurred when reading the file specified in the type command. Check the file name and path. For more information, see the type command can diagnose trouble.

READER STOPPING

PIP. Reader operation interrupted.

Record Too Long

PIP.	PIP cannot process a record longer than 128 bytes.	
------	--	--

Requires CP/M 2.0 or later

XSUB. XSUB requires the facilities of CP/M 2.0 or newer version.



Message	Meaning
---------	---------

PIP. This version of PIP requires the facilities of CP/M 2.0 or newer version.

```
PIP. The string argument to an S
parameter cannot be found in the source
file.
```

```
SYSGEN. SYSGEN cannot use your CP/M
source file.
```

ASM. When you assemble a file, you cannot use the wildcard characters \* and ? in the

ASM. The assembler cannot understand the information in the file containing the assembly-language program. Portions of another file might have been written over your assembly-language file, or information was not properly saved on the disk. Use the TYPE command to locate the error. Assembly-language files contain the letters, symbols, and numbers that appear on your keyboard. If your screen displays unrecognizable output or behaves strangely, you have found where computer instructions have crept into your file.

MOVCPM. The MOVCPM utility is being used with the wrong CP/M system.



CP/M Operating System Manual	I CP/M Error Messages
------------------------------	-----------------------

Table I-1. (continued)

[illegible]

```
"SYSTEM" FILE NOT ACCESSIBLE
```

```
You tried to access a file set to SYS with
the STAT command.
```

```

**      TOO MANY FILES      **

```

```
STAT. There is not enough memory for STAT
to sort the files specified, or more than
512 files were specified.
```

```
UNEXPECTED END OF HEX FILE--{filespec}
```

P1P. An end-of-file was encountered prior to a termination HEX record. The HEX file without a termination record should be corrected, probably by recreating the HEX file

Unrecognized Destination	Recognized Destination

PIP.	Check	command	line	for	valid
	destination				

```
Use: STAT d:=RO
```

```
STAT.  An invalid STAT drive command was
given.  The only valid drive assignment in
STAT is STAT d:=RO.
```

```
VERIFY ERROR:--{filespec}
```

PTP. When copying with the M option, PTP found a difference when rereading the data just written and comparing it to the data in its memory buffer. Usually this indicates a failure of either the destination disk or drive.

WRONG CP/M VERSION (REQUIRES 2.0)

XSUB	ACTIVE
------	--------

SUBMIT XSUB has been invoked

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------



[illegible]

XSUB	ALREADY	PRESENT
------	---------	---------

SUBMIT. XSUB is already active in memory.

Your input?

If CP/M cannot find the command you specified, it returns the command name you entered followed by a question mark. Check that you have typed the command line correctly, or that the command you requested exists as a .COM file on the default or specified disk.

E	n	d	o	f	A	p	p	e	n	d	i	x	I
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[illegible]



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